

*Application Note 103*  
*6x86MX BIOS*  
*WRITER'S GUIDE*



---

## Table of Contents

1.0	Introduction	
1.1	Scope	4
1.2	Cyrix Configuration Registers	5
1.3	Summary of 6x86MX and 6x86 Differences	6
2.0	Cache Unit	
3.0	Cyrix 6x86MX CPU Detection	
3.1	Detecting the Cyrix 6x86MX - Method 1	8
3.2	Detecting the Cyrix 6x86MX - Method 2	10
3.3	EDX Value Following Reset	12
3.4	Determining 6x86MX Operating Frequency	13
3.5	Device Name	14
4.0	6x86MX Configuration Register Index Assignments	
4.1	Accessing a Configuration Register	15
4.2	6x86MX Configuration Register Index Assignments	15
4.3	Configuration Control Registers (CCR0-6)	17
4.4	Address Region Registers (ARR0-7)	25
4.5	Region Control Registers (RCR0-7)	27
5.0	Recommended 6x86MX Configuration Register Settings	
5.1	PC Memory Model	32
5.2	General Recommendations	34
5.3	Recommended Bit Settings	36
6.0	Model Specific Registers	
6.1	Time Stamp Counter	43
6.2	Performance Monitoring	43
6.3	Performance Monitoring Counters 1 and 2	44
6.4	Counter Event Counter Register	44
6.5	PM Pin Control	44

---

7.0	Programming Model Differences	
7.1	Instruction Set	50
7.2	Configuring Internal 6x86MX Features	50
7.3	INVD and WBINVD Instructions	50
7.4	Control Register 0 (CR0) CD and NW Bits	51
Appendixes		
Appendix A	-Sample Code: Detecting a Cyrix CPU	52
Appendix B	-Sample Code: Determining CPU MHz	54
Appendix C	-Example CPU Type and Frequency Detection Program	57
Appendix D	-Sample Code: Programming 6x86MX Configuration Registers	59
Appendix E	-Sample Code: Controlling the L1 Cache	60
Appendix F	-Example Configuration Register Settings	61
Appendix G	-Sample Code: Detecting L2 Cache Burst Mode	62

---

## *1 Introduction*

---

### *1.1 Scope*

This document is intended for 6x86MX system BIOS writers. It is not a stand-alone document, but a supplement to other Cyrix documentation including the 6x86MX Data Book, and Cyrix SMM Programmer's Guide. This document highlights the programming differences between the 6x86 and the 6x86MX. Recommendations for 6x86MX detection and configuration register settings are included.

The recommended settings are optimized for performance and compatibility in Windows95 or Windows NT, Plug and Play (Pnp), PCI-based system. Performance optimization, CPU detection, chipset initialization, memory discovery, I/O recovery time, and other functions are described in detail.

---

## *1.2 Cyrix Configuration Registers*

The 6x86MX uses on-chip configuration registers to control the on-chip cache, system management mode (SMM), device identification, and other 6x86MX specific features. The on-chip registers are used to activate advanced performance features. These performance features may be enabled “globally” in some cases, or by a user-defined address region. The flexible configuration of the 6x86MX is intended to fit a wide variety of systems.

### **The Importance of Non-Cacheable Regions**

The 6x86MX has eight internal user-defined Address Region Registers. Among other attributes, the regions define cacheability of the address regions. Using this cacheability information, the 6x86MX is able to implement high performance features, that would otherwise not be possible. A non-cacheable region implies that read sourcing from the write buffers, data forwarding, data bypassing, speculative reads, and fill buffer streaming are disabled for memory accesses within that region. Additionally, strong cycle ordering is also enforced. Although negating KEN# during a memory access on the bus prevents a cache line fill, it does not fully disable these performance features. In other words, negating KEN# is NOT equivalent to establishing a non-cacheable region in the 6x86MX.

---

### *1.3 Summary of 6x86MX and 6x86 Differences*

The differences between the 6x86MX CPU and the 6x86 CPU are listed in the table below.

	<b>6x86MX</b>	<b>6x86</b>	<b>NOTES</b>
L1 Cache Size	64 KBytes	16 KBytes	Section 2
CPUID (Bit 7 of CCR4)	Reset Value = 1	Reset Value = 0	Section 3 Section 4
Family Code	06h	05h	Figure 3-1
EDX	Reset Value = 06 + DIR0	Reset Value = 05 + DIR0	Section 3
Time Stamp Counter	Yes	No	BIOS Writer's Guide Revision 1.2
DIR0 (Register Index = FEh)	5xh	3xh	Section 3
DTE_EN (Bit 4 of CCR4)	Reserved	If = 1, the DTE cache is enabled.	Section 4
SLOP (Bit 1 of CCR5)	Reserved	If =1, the LOOP instruction is slowed down.	Section 4
LBR1 (Bit 4 of CCR5)	Reserved	If =1, LBA# pin is asserted for all accesses to the 640KBytes - 1MByte address region.	Section 4
WVO (Bit 1 of RCRx)	Reserved	If = 1, weak write ordering is enabled for the corresponding region.	Section 4

**Summary of 6x86MX and 6x86 Differences**

---

## *2. Cache Unit*

The cache size of the 6x86MX has been increased to 64 KByte. This is four times larger than the 16 KByte cache of the 6x86. The cache is configured the same way as the 6x86: 4-way set associative, and 32 Byte lines.

---

## *3. Cyrix 6x86MX CPU Detection*

Two methods for detecting the Cyrix 6x86MX CPU are described in Sections 3.1 and 3.2.

Cyrix does not recommend other detection algorithms using the value of EDX following reset, and other signature methods of determining if the CPU is an 8086, 80286, 80386, or 80486.

---

### *3.1 Detecting the Cyrix 6x86MX - Method 1*

This method for detecting the presence of an 6x86MX microprocessor during BIOS POST is a two step process. First, a Cyrix brand CPU must be detected. Second, the CPU's Device Identification Registers (DIRs) provide the CPU model and stepping information.

---

#### *3.1.1 Cyrix CPU Detection*

Detection of a Cyrix brand CPU is implemented by checking the state of the undefined flags following execution of the divide instruction which divides 5 by 2 (5÷2). The undefined flags in a Cyrix microprocessor remain unchanged following the divide. Alternate CPUs modify some of the undefined flags. Using operands other than 5 and 2 may prevent the algorithm from working correctly. Appendix A contains sample code for detecting a Cyrix CPU using this method.

---

#### *3.1.2 Detecting CPU Type and Stepping*

Once a Cyrix brand CPU is detected, the model and stepping of the CPU can be determined. All Cyrix CPUs contain Device Identification Registers (DIRs) that exist as part of the configuration registers. The DIRs for all Cyrix CPUs exist at configuration register indexes 0FEh and 0FFh. The table below specifies the contents of the 6x86MX DIRs.

DIR0 bits [7:4] = 5h indicate an 6x86MX CPU is present, DIR0 bits [3:0] indicate the core-to-bus clock ratio, and DIR1 contains stepping information. Clock ratio information is provided to assist calculations in determining bus frequency once the CPU's core frequency has been calculated. Proper bus speed settings are critical to overall system performance.

DEVICE	CORE/BUS CLOCK RATIO	DIR0 (DEVICE ID)	DIR1 (REV ID)
6x86MX	2/1 (default) 2.5/1 3/1 3.5/1	51h or 59h 52h or 5Ah 53h or 5Bh 54h or 5Ch	TBD

**Cyrix Device Identification Registers**

---

### 3.2 *Detecting the Cyrix 6x86MX - Method 2*

Unlike the 6x86, the CPUID instruction is enabled following reset. It can be disabled by clearing the CPUID bit in configuration register CCR4. It is recommended that all BIOS vendors include a CPUID enable/disable field in the CMOS setup to allow the end-user to disable the CPUID instruction.

The CPUID instruction, opcode 0FA2h, provides information indicating Cyrix as the vendor and the family, model, stepping, and CPU features. The EAX register provides the input value for the CPUID instruction. The EAX register is loaded with a value to indicate what information should be returned by the instruction.

Following execution of the CPUID instruction with an input value of "0" in EAX, the EAX, EBX, ECX and EDX registers contain the information shown in Figure 3-1. EAX contains the highest input value understood by the CPUID instruction, which for the 6x86MX is "1". EBX, ECX and EDX contain the vendor identification string "CyrixInstead".

Following execution of the CPUID instruction with an input value of "1" loaded in EAX, EAX[15:0] will contain the value of 06xxh. EDX [31-0] will contain the value 0080A135h.

```

switch (EAX)
{
  case (0):
    EAX := 1
    EBX := 69 72 79 43/* 'i' 'r' 'y' 'C' */
    EDX := 73 6e 49 78/* 's' 'n' 'I' 'x' */
    ECX := 64 61 65 74/* 'd' 'a' 'e' 't' */
    break

  case (1):
    EAX[7:0] := 00h
    EAX[15:8] := 06h
    EDX[0] := 1 /* 1=FPU Built In */
    EDX[1] := 0 /* 0=No V86 enhancements */
    EDX[2] := 1 /* 1=I/O breakpoints */
    EDX[3] := 0 /* 0=No page size extensions */
    EDX[4] := 1 /* 1=Time Stamp Counter */
    EDX[5] := 1 /* 1=RDMSR and WRMSR */
    EDX[6] := 0 /* 0=No physical address extensions */
    EDX[7] := 0 /* 0=No machine check exception */
    EDX[8] := 1 /* 1=CMOV, FCMOV, FCOMI instructions */
    EDX[9] := 0 /* 0=No APIC */
    EDX[11-10] := 0 /* Undefined */
    EDX[12] := 0 /* 0=No memory type range registers */
    EDX[13] := 1 /* 1=PTE global bit */
    EDX[14] := 0 /* 0=No machine check architecture */
    EDX[15] := 1 /* 1=CMOV, FCMOV, FCOMI instructions */
    EDX[22-16] := 0 /* Undefined */
    EDX[23] := 1 /* 1=MMX instructions */
    EDX[31-24] := 00h /* "documentation error was: 0080h" */

    break

  default:
    EAX, EBX, ECX, EDX : Undefined
}

```

### Information Returned by CPUID Instruction

---

### *3.3 EDX Value Following Reset*

Some CPU detection algorithms may use the value of the CPU's EDX register following reset. The 6x86MX's EDX register contains the data shown below following a reset initiated using the RESET pin:

EDX[31:16] = undefined  
EDX[15:8] = 06h  
EDX[7:0] = DIR0

Refer to the table on the previous page for DIR0 values. The value in EDX does not identify the vendor of the CPU. Therefore, EDX alone cannot be used to determine if a Cyrix CPU is present. However, BIOS should preserve the contents of EDX so that applications can use the EDX value when performing a user-defined shutdown, e.g. a reset performed with data 0Ah in the Shutdown Status byte (Index 0Fh) of the CMOS RAM map.

---

### 3.4 *Determining 6x86MX Operating Frequency*

Determining the operating frequency of the CPU is normally required for correct initialization of the system logic. Typically, a software timing loop with known instruction clock counts is timed using legacy hardware (the 8254 timer/counter circuits) within the PC. Once the operating frequency of the 6x86MX's core is known, DIR0 bits (2:0) can be examined to calculate the bus operating frequency.

---

#### 3.4.1 *Instruction Count Method*

Careful selection of instructions and operands must be used to replicate the exact clock counts detailed in the Instruction Set Summary found in the 6x86MX Data Book. An example code sequence for determining the 6x86MX's operating frequency is detailed in Appendix B and Appendix C. This code sequence is identical to the recommended sequence for the 6x86. The core loop uses a series of five IDIV instructions within a LOOP instruction. IDIV was chosen because it is an exclusive instruction meaning that it executes in the 6x86MX x-pipeline with no other instruction in the y-pipeline. This allows for more predictable execution times as compared to using non-exclusive instructions.

The 6x86MX instruction clock count for IDIV varies from 17 to 45 clocks for a doubleword divide depending on the value of the operands. The code example in the appendices uses "0" divided by "1" which takes only 17 clocks to complete. The LOOP instruction clock count is 1. Therefore, the overall clock count for the inner loop in this example is 86 clocks.

---

#### 3.4.2 *Time Stamp Counter Method*

On the 6x86MX, the Time Stamp Counter (TSC) can be used as an alternative method for obtaining an exact core clock count during the software timing loop.

The Time Stamp Counter is a 64-bit counter that counts internal CPU clock cycles since the last reset. The value can be read any time via the RDTSC instruction, opcode OF31h. The RDTSC instruction loads the contents of the TSC into EDX:EAX. The use of the RDTSC instruction is restricted by the Time Stamp Disable, (TSD) flag in CR4. When the TSD flag is 0, the RDTSC instruction can be executed at any privilege level. When the TSD flag is 1, the RDTSC instruction can only be executed at privilege level 0.

---

**Device Name**

The exact core count during the software timing loop can be determined by computing the difference of the Time Stamp Counter at start of the loop and the end of the loop.

---

### 3.5 *Device Name*

The correspondence between core frequency, bus frequency and performance rating is shown in the table below. The device name in table below should be used by the BIOS for display during boot-up and in BIOS setup screens or utilities.

DEVICE NAME AND PART NUMBER	CLOCK MULTIPLIER	FREQUENCY (MHz)	
		BUS	INTERNAL
6x86MX - PR166GP	2.5	60	150
6x86MX - PR200GP	2.5	66	166
6x86MX - PR233GP	2.5	75	188
6x86MX - PR233GP	3.0	66	200
6x86MX - PR266GP	3.0	75	225
6x86MX - PR266GP	3.5	66	233

**Cyrix 6x86MX™ Part Numbers**

---

## *4. 6x86MX Configuration Register Index Assignments*

On-chip configuration registers are used to control the on-chip cache, system management mode and other 6x86MX unique features.

---

### *4.1 Accessing a Configuration Register*

Access to the configuration registers is achieved by writing the index of the register to I/O port 22h. I/O port 23h is then used for data transfer. Each I/O port 23h data transfer must be preceded by an I/O port 22h register index selection, otherwise the second and later I/O port 23h operations are directed off-chip and produce external I/O cycles. Reads of I/O port 22h are always directed off-chip. Appendix D contains example code for accessing the 6x86MX configuration registers.

---

### *4.2 6x86MX Configuration Register Index Assignments*

The table on the following page lists the 6x86MX configuration register index assignments. After reset, configuration registers with indexes C0-CFh and FC-FFh are accessible. In order to prevent potential conflicts with other devices which may use ports 22 and 23h to access their registers, the remaining registers (indexes 00-BFh, D0-FBh) are accessible only if the MAPEN(3-0) bits in CCR3 are set to 1h. With MAPEN(3-0) set to 1h, any access to an index in the 00-FFh range does not create external I/O bus cycles. Registers with indexes C0-CFh, FC-FFh are accessible regardless of the state of the MAPEN bits. If the register index number is outside the C0-CFh or FE-FFh ranges, and MAPEN is set to 0h, external I/O bus cycles occur. The table on the next page lists the MAPEN values required to access each 6x86MX configuration register. The configuration registers are described in more detail in the following sections.

---

**6x86MX Configuration Register Index Assignments**

REGISTER INDEX	REGISTER NAME	ACRONYM	WIDTH (BITS)	MAPEN(3-0)
00h-BFh	Reserved	—	—	—
C0h	Configuration Control 0	CCR0	8	Don't care
C1h	Configuration Control 1	CCR1	8	Don't care
C2h	Configuration Control 2	CCR2	8	Don't care
C3h	Configuration Control 3	CCR3	8	Don't care
C4h-C6h	Address Region 0	ARR0	24	Don't care
C7h-C9h	Address Region 1	ARR1	24	Don't care
CAh-CCCh	Address Region 2	ARR2	24	Don't care
CDh-CFh	Address Region 3	ARR3	24	Don't care
D0h-D2h	Address Region 4	ARR4	24	1h
D3h-D5h	Address Region 5	ARR5	24	1h
D6h-D8h	Address Region 6	ARR6	24	1h
D9h-DBh	Address Region 7	ARR7	24	1h
DCh	Region Configuration 0	RCR0	8	1h
DDh	Region Configuration 1	RCR1	8	1h
DEh	Region Configuration 2	RCR2	8	1h
DFh	Region Configuration 3	RCR3	8	1h
E0h	Region Configuration 4	RCR4	8	1h
E1h	Region Configuration 5	RCR5	8	1h
E2h	Region Configuration 6	RCR6	8	1h
E3h	Region Configuration 7	RCR7	8	1h
E4h-E7h	Reserved	—	—	—
E8h	Configuration Control 4	CCR4	8	1h
E9h	Configuration Control 5	CCR5	8	1h
EAh	Configuration Control 6	CCR6	8	1h
EBh-FAh	Reserved	—	—	—
FBh	Device Identification 2	DIR2	8	1h
FCh	Device Identification 3	DIR3	8	1h
FDh	Device Identification 4	DIR4	8	1h
FEh	Device Identification 0	DIR0	8	Don't care
FFh	Device Identification 1	DIR1	8	Don't care

**Configuration Register Index Assignments**

The 6x86MX configuration registers can be grouped into four areas:

- Configuration Control Registers (CCRs)
- Address Region Registers (ARRs)
- Region Control Registers (RCRs)
- Device Identification Registers (DIRs)

CCR bits independently control 6x86MX features. ARR and RCR define regions of memory with specific attributes. DIRs are used for CPU detection as discussed earlier in Chapter 3. All bits in the configuration registers are initialized to zero following reset unless specified otherwise. The appropriate configuration register bit settings vary depending on system design. Optimal settings recommended for a typical PC environment are discussed in Chapter 5.

---

### *4.3 Configuration Control Registers (CCR0-6)*

There are seven CCRs in the 6x86MX which control the cache, power management and other unique features. The following paragraphs describe the CCRs and associated bit definitions in detail.

---

**Configuration Control Registers (CCR0-6)**

---

**4.3.1 Configuration Control Register 0 (CCR0)**

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
Reserved	Reserved	Reserved	Reserved	Reserved	Reserved	NC1	Reserved

**Configuration Control Register 0 (CCR0)**

BIT NAME	BIT No.	DESCRIPTION
NC1	1	If = 1, designates 640KBytes -1MByte address region as non-cacheable. If = 0, designates 640KBytes -1MByte address region as cacheable.

**CCR0 Bit Definitions**

---

4.3.2 Configuration Control Register 1 (CCR1)

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
SM3	Reserved	Reserved	NO_LOCK	Reserved	SMAC	USE_SMI	Reserved

**Configuration Control Register 1 (CCR1)**

BIT NAME	BIT No.	DESCRIPTION
SM3	7	If = 1, designates Address Region Register 3 as SMM address space.
NO_LOCK	4	If = 1, all bus cycles are issued with the LOCK# pin negated except page table accesses and interrupt acknowledge cycles. Interrupt acknowledge cycles are executed as locked cycles even though LOCK# is negated. With NO_LOCK set, previously non-cacheable locked cycles are executed as unlocked cycles and therefore, may be cached. This results in higher CPU performance. See the section on Region Configuration Registers (RCR) for more information on eliminating locked CPU bus cycles only in specific address regions.
SMAC	2	If = 1, any access to addresses within the SMM address space access system management memory instead of main memory. SMI# input is ignored while SMAC is set. Setting SMAC=1 allows access to SMM memory without entering SMM. This is useful for initializing or testing SMM memory.
USE_SMI	1	If = 1, SMI# and SMIACT# pins are enabled. If = 0, SMI# pin is ignored and SMIACT# pin is driven inactive.

**CCR1 Bit Definitions**

---

**Configuration Control Registers (CCR0-6)**

---

**4.3.3 Configuration Control Register 2 (CCR2)**

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
USE_SUSP	Reserved	Reserved	WPR1	SUSP_HLT	LOCK_NW	SADS	Reserved

**Configuration Control Register 2 (CCR2)**

BIT NAME	BIT No.	DESCRIPTION
USE_SUSP	7	If = 1, SUSP# and SUSPA# pins are enabled. If = 0, SUSP# pin is ignored and SUSPA# pin floats. These pins should only be enabled if the external system logic (chipset) supports them.
WPR1	4	If = 1, designates that any cacheable accesses in the 640 KBytes-1MByte address region are write-protected. With WPR1=1, any attempted write to this range will not update the internal cache.
SUSP_HLT	3	If = 1, execution of the HLT instruction causes the CPU to enter low power suspend mode. This bit should be used with caution since the CPU must recognize and service an INTR, NMI or SMI to exit the "HLT initiated" suspend mode.
LOCK_NW	2	If = 1, the NW bit in CR0 becomes read only and the CPU ignores any writes to this bit.
SADS	1	If = 1, the CPU inserts an idle cycle following sampling of BRDY# and prior to asserting ADS#.

**CCR2 Bit Definitions**

4.3.4 Configuration Control Register 3 (CCR3)

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
MAPEN				Reserved	LINBRST	NMI_EN	SML_LOCK

**Configuration Control Register 3 (CCR3)**

BIT NAME	BIT No.	DESCRIPTION
MAPEN	7-4	If set to 0001 binary (1h), all configuration registers are accessible. If set to 0000, only configuration registers with indices C0-CFh, FEh and FFh are accessible.
LINBRST	2	If = 1, the 6x86MX will use a linear address sequence when performing burst cycles. If = 0, the 6x86MX will use a "1+4" address sequence when performing burst cycles. The "1+4" address sequence is compatible with the Pentium's burst address sequence.
NMI_EN	1	If = 1, NMI interrupt is recognized while in SMM. This bit should only be set while in SMM, after the appropriate NMI interrupt service routine has been setup.
SML_LOCK	0	If = 1, the CPU prevents modification of the following SMM configuration bits, except when operating in an SMM service routine: CCR1 USE_SMI, SMAC, SM3 CCR3 NMI_EN ARR3 Starting address and block size. Once set, the SML_LOCK bit can only be cleared by asserting the RESET pin.

**CCR3 Bit Definitions**

---

**Configuration Control Registers (CCR0-6)**

---

**4.3.5 Configuration Control Register 4 (CCR4)**

The 6x86 DTE cache has been eliminated on the 6x86MX. Therefore, bit 4 of CCR4 is a reserved bit.

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
CPUID	Reserved	Reserved	Reserved	Reserved	IORT		

**Configuration Control Register 4 (CCR4)**

BIT NAME	BIT No.	DESCRIPTION
CPUID	7	If = 1, bit 21 of the EFLAG register is write/readable and the CPUID instruction will execute normally. If = 0, bit 21 of the EFLAG register is not write/readable and the CPUID instruction is an invalid opcode.
IORT	2-0	Specifies the minimum number of bus clocks between I/O accesses (I/O recovery time). The delay time is the minimum time from the end of one I/O cycle to the beginning of the next (i.e. BRDY# to ADS# time). 0h = 1 clock 1h = 2 clocks 2h = 4 clocks 3h = 8 clocks 4h = 16 clocks 5h = 32 clocks (default value after RESET) 6h = 64 clocks 7h = no delay

**CCR4 Bit Definitions**

---

#### 4.3.6 Configuration Control Register 5 (CCR5)

The 6x86 Slow Loop Instruction and Local Bus Access features have been eliminated in the 6x86MX. Therefore, bits 4 and 1 of CCR5 are reserved bits on the 6x86MX.

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
Reserved	Reserved	ARREN	Reserved	Reserved	Reserved	Reserved	WT_ALLOC

#### Configuration Control Register 5 (CCR5)

BIT NAME	BIT No.	DESCRIPTION
ARREN	5	If = 1, enables all Address Region Registers (ARRs). If clear, disables the ARR registers. If SM3 is set, ARR3 is enabled regardless of the ARREN setting.
WT_ALLOC	0	If = 1, new cache lines are allocated for both read misses and write misses. If = 0, new cache lines are only allocated on read misses.

#### CCR5 Bit Definitions

---

**Configuration Control Registers (CCR0-6)**

---

**4.3.7 Configuration Control Register 6 (CCR6)**

Configuration Control Register 6 has been added to the 6x86MX.

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
Reserved	N	Reserve	Reserved	Reserved	Reserved	WP_ARR3	SMM_MODE

**Configuration Control Register 6 (CCR6)**

BIT NAME	BIT No.	DESCRIPTION
N	6	Nested SMI Enable bit: If operating in Cyrix enhanced SMM mode and: If = 1: Enables nesting of SMI's If = 0: Disable nesting of SMI's. This bit is automatically CLEARED upon entry to every SMM routine and is SET upon every SMM routine and is SET upon every RSM. Therefore enabling/disabling of nested SMI can only be done while operating in SMM mode.
WP_ARR3	1	If = 1: Memory region defined by ARR3 is write protected when operating outside of SMM mode. If = 0: Disable write protection for memory region defined by ARR3. Reset State = 0.
SMM_MODE	0	If = 1: Enables Cyrix Enhanced SMM mode. If = 0: Disables Cyrix Enhanced SMM mode.

**CCR6 Bit Definitions**

#### 4.4 Address Region Registers (ARR0-7)

The Address Region Registers (ARRs) are used to define up to eight memory address regions. Each ARR has three 8-bit registers associated with it which define the region starting address and block size. The Table “ARRx Index Assignments” below shows the general format for each ARR and lists the index assignments for the ARR’s starting address and block size.

The region starting address is defined by the upper 12 bits of the physical address. The region size is defined by the BSIZE(3-0) bits as shown in the Table “BSIZE (3-0) Bit Definitions” on the next page. The BIOS and/or its utilities should allow definition of all ARR. There is one restriction when defining the address regions using the ARR. The region starting address must be on a block size boundary. For example, a 128KByte block is allowed to have a starting address of 0KBytes, 128KBytes, 256KBytes, and so on.

ADDRESS REGION REGISTER	STARTING ADDRESS			REGION BLOCK SIZE
	A31-A24	A23-A16	A15-A12	BSIZE(3-0)
	BITS (7-0)	BITS (7-0)	BITS (7-4)	BITS (3-0)
ARR0	C4h	C5h		C6h
ARR1	C7h	C8h		C9h
ARR2	CAh	CBh		CCh
ARR3	CDh	CEh		CFh
ARR4	D0h	D1h		D2h
ARR5	D3h	D4h		D5h
ARR6	D6h	D7h		D8h
ARR7	D9h	DAh		DBh

**ARRx Index Assignments**

---

**Address Region Registers (ARR0-7)**

<b>BSIZE(3-0)</b>	<b>ARR(0-6) REGION SIZE</b>	<b>ARR7 REGION SIZE</b>
0h	Disabled	Disabled
1h	4 KBytes	256 KBytes
2h	8 KBytes	512 KBytes
3h	16 KBytes	1 MByte
4h	32 KBytes	2 MBytes
5h	64 KBytes	4 MBytes
6h	128 KBytes	8 MBytes
7h	256 KBytes	16 MBytes
8h	512 KBytes	32 MBytes
9h	1 MByte	64 MBytes
Ah	2 MBytes	128 MBytes
Bh	4 MBytes	256 MBytes
Ch	8 MBytes	512 MBytes
Dh	16 MBytes	1 GBytes
Eh	32 MBytes	2 GBytes
Fh	4 GBytes	4 GBytes

**BSIZE (3-0) Bit Definitions**

---

**Region Control Registers (RCR0-7)**

---

**4.5 Region Control Registers (RCR0-7)**

The RCRs are used to define attributes, or characteristics, for each of the regions defined by the ARR. Each ARR has a corresponding RCR with the general format shown below.

New to the 6x86MX is the Invert Region feature. This feature is controlled by the INV\_RGN bit of the Region Control Registers.

If the INV\_RGN bit is set, the controls specified in the RCR (RCD, WT, WG, WL) will be applied to all memory addresses outside the region specified in the corresponding ARR.

If the INV\_RGN bit is cleared, the 6x86MX functions identically to the 6x86 (the controls specified in the RCR will be applied to all memory addresses inside the region specified by the corresponding ARR).

The INV\_RGN bit is defined for RCR(0-6) only. 6x86 Weak Write Ordering and Local Bus Access features have been eliminated on the 6x86MX. Therefore, bit 5 and bit 1 are reserved bits for the 6x86MX.

BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1	BIT 0
Reserved	INV_RGN	Reserved	WT	WG	WL	Reserved	RCD/RCE

**RCR Bit Definitions**

Note: RCD is defined for RCR0-RCR6. RCE is defined for RCR7 only.

BIT NAME	BIT No.	DESCRIPTION
RCD	0	Applicable to RCR(0-6) only. If set, the address region specified by the corresponding ARR is non-cacheable.
RCE	0	Applicable to RCR7 only. If set, the address region specified by ARR7 is cacheable and implies that address space outside of the region specified by ARR7 is non-cacheable.
WL	2	If set, weak locking is enabled for the corresponding region.
WG	3	If set, write gathering is enabled for the corresponding region.
WT	4	If set, write through caching is enabled for the corresponding region.
INV_RGN	6	Applicable to RCR(0-6) only. If set, apply controls specified in RCR to all memory addresses outside the region specified in the corresponding ARR.

**RCR Bit Definitions**

---

#### *4.5.1 Detailed Description of RCR Attributes*

##### **Region Cache Disable (RCD)**

Setting RCD=1 defines the corresponding address region as non-cacheable. RCD prevents caching of any access within the specified region. Additionally, RCD implies that high performance features are disabled for accesses within the specified address region. Bus cycles issued to memory addresses within the specified region are single cycles with the CACHE# pin negated. If KEN# is asserted for a memory access within a region defined non-cacheable by RCD, the access is not cached.

##### **Region Cache Enable (RCE)**

Setting RCE=1 defines the corresponding address region as cacheable. RCE is applicable to ARR7 only. RCE in combination with ARR7, is intended to define the Main Memory Region. All memory outside ARR7 is non-cacheable when RCE is set. This is intended to define all unused memory space as non-cacheable. If KEN# is negated for an access within a region defined cacheable by RCE, the access is not cached.

##### **Weak Locking (WL)**

Setting WL=1 enables weak locking for the corresponding address region. With WL enabled, all bus cycles are issued with the LOCK# pin negated except for page table accesses and interrupt acknowledge cycles. WL negates bus locking so that previously non-cacheable cycles can be cached. Typically, XCHG instructions, instructions preceded by the LOCK prefix, and descriptor table accesses are locked cycles. Setting WL allows the data for these cycles to be cached.

Weak Locking (WL) implements the same function as NO\_LOCK except that NO\_LOCK is a global enable. The NO\_LOCK bit of CCR1 enables weak locking for the entire address space, whereas the WL bit enables weak locking only for specific address regions.

##### **Write Gathering (WG)**

Setting WG=1 enables write gathering for the corresponding address region. With WG enabled, multiple byte, word or dword writes to sequential addresses that would normally occur as individual write cycles are combined and issued as a sin-

gle write cycle. WG improves bus utilization and should be used for memory regions that are not sensitive to the “gathering.” WG can be enabled for both cacheable and non-cacheable regions.

### Write Through (WT)

Setting WT=1 defines the corresponding address region as write-through instead of write-back. Any system ROM that is allowed to be cached by the processor should be defined as write-through.

---

#### 4.5.2 *Attributes for Accesses Outside Defined Regions*

If an address is accessed that is not in a region defined by the ARR<sub>n</sub> and ARR<sub>7</sub> is defined with RCE=1, the following conditions apply:

- The memory access is not cached regardless of the state of KEN#.
- Writes are not gathered.
- Strong locking occurs.
- Strong write ordering occurs.

---

#### 4.5.3 *Attributes for Accesses in Overlapped Regions*

If two defined address regions overlap (including NC1 and LBR1) and conflicting attributes are specified, the following attributes take precedence:

- Write-back is disabled.
- Writes are not gathered.
- Strong locking occurs.
- Strong write ordering occurs.
- The overlapping regions are non-cacheable.

---

**Region Control Registers (RCR0-7)**

Since the CCR0 bit NC1 affects cacheability, a potential exists for conflict with the ARR7 main memory region which also affects cacheability. This overlap in address regions causes a conflict in cacheability. In this case, NC1 takes precedence over the ARR7/RCE setting because non-cacheability always takes precedence. For example, for the following settings:

- NC1=1
- ARR7 = 0-16 MBytes
- RCR7 bit RCE = 1

The 6x86MX caches accesses as shown in the table below.

ADDRESS REGION	CACHEABLE	COMMENTS
0 to 640 KBytes	Yes	ARR7/RCE setting.
640 KBytes- 1 MByte	No	NC1 takes precedence over ARR7/RCE setting.
1 MByte - 16 MBytes	Yes	ARR7/RCE setting.
16 MBytes - 4 GBytes	No	Default setting.

**Cacheability for Example 1**

---

#### 4.5.4 *Attributes for Accesses with Conflicting Signal Pin Inputs*

The characteristics of the regions defined by the ARR and the RCR may also conflict with indications by hardware signals (i.e., KEN#, WB/WT#). The following paragraphs describe how conflicts between register settings and hardware indicators are resolved.

##### Non-cacheable Regions and KEN#

Regions which have been defined as non-cacheable (RCD=1) by the ARR and RCR may conflict with the assertion of the KEN# input. If KEN# is asserted for an access to a region defined as non-cacheable, the access is not cached. Regions defined as non-cacheable by the ARR and RCR take precedence over KEN#. The NC1 bit also takes precedence over the KEN# pin. If NC1 is set, any access to the 640 KByte-1 MByte address region with KEN# asserted is not cached.

##### Write-Through Regions and WB/WT#

Regions which have been defined as write-through (WT=1) may conflict with the state of the WB/WT# input to the 6x86MX. Regions defined as write-through by the ARR and RCR remain write-through even if WB/WT# is asserted during accesses to these regions. The WT bit in the RCR takes precedence over the state of the WB/WT# pin in cases of conflict.

## 5. Recommended 6x86MX Configuration Register Settings

### 5.1 PC Memory Model

The table below defines the allowable attributes for a typical PC memory model. Actual recommended configuration register settings for a typical PC system are listed in Appendix F.

ADDRESS SPACE	ADDRESS RANGE	CACHEABLE	WEAK LOCKS	WRITE GATHERED	WRITE-THROUGH	NOTES
DOS Area	0-9 FFFFh	Yes	No	Yes	No	
Video Buffer	A 0000-B FFFFh	No	No	Yes	No	Note 1
Video ROM	C 0000-C 7FFFh	Yes	No	No	Yes	Note 2
Expansion Card/ROM Area	C 8000h-D FFFFh	No	No	No	No	
System ROM	E 0000h-F FFFFh	Yes	No	No	Yes	Note 2
Extended Memory	10 0000h- Top of Main Memory	Yes	No	Yes	No	
Unused/PCI MMIO	Top of Main Memory- FFFF FFFFh	No	No	No	No	Note 3

#### PC Memory Model

##### Notes 1: Video Buffer Area

A non-cacheable region must be used to enforce strong cycle ordering in this area and to prevent caching of Video RAM. The Video RAM area is sensitive to bus cycle ordering. The VGA controller can perform logical operations which depend on strong cycle ordering (found in Windows 3.1 code). To guarantee that the 6x86MX performs strong cycle ordering, a non-cacheable area must be established to cover the Video RAM area.

Video performance is greatly enhanced by gathering writes to Video RAM. For example, video performance benchmarks have been found to use REP STOSW instructions that would normally execute as a series of sequential 16-bit write cycles. With WG enabled, groups of four 16-bit write cycles are reduced to a single 64-bit write cycle.

Note 2: Video ROM and System ROM

Caching of the Video and System ROM areas is permitted, but is normally non-cacheable because NC1 is set. If these areas are cached, they must be cached as write-through regions. 6x86MX system benchmarking in a Windows environment has shown no benefit to caching these ROM areas. Therefore, it is recommended that these areas be set as non-cacheable using the NC1 bit in CCR0.

Note 3: Top of Main Memory-FFFF FFFFh (Unused/PCI Memory Space)

Unused/PCI Memory Space immediately above physical main memory must be defined as non-cacheable to ensure proper operation of memory sizing software routines and to guarantee strong cycle ordering. Memory discovery routines must occur with cache disabled to prevent read sourcing from the write buffers. Also, PCI memory mapped I/O cards that may exist in this address region may contain control registers or FIFOs that depend on strong cycle ordering.

The appropriate non-cacheable region must be established using ARR7. For example, if 32 MBytes (000 0000h-1FF FFFFh) are installed in the system, a non-cacheable region must begin at the 32 MByte boundary (200 0000h) and extend through the top of the address space (FFFF FFFFh). This is accomplished by using ARR7 (Base = 0000 0000h, BSize = 32 MBytes) in combination with RCE=1.

---

## 5.2 *General Recommendations*

---

### 5.2.1 *Main Memory*

Memory discovery routines should always be executed with the L1 cache disabled. By default, L1 caching is globally disabled following reset because the CD bit in Control Register 0 (CR0) is set. Always ensure the L1 cache is disabled by setting the CD bit in CR0 or by programming an ARR to “4 GByte cache disabled” before executing the memory discovery routine. Once BIOS completes memory discovery, ARR7 should be programmed with a base address of 000 0000h and with a “Size” equal to the amount of main memory that was detected.

The intent of ARR7 is to define a cacheable region for main memory and simultaneously define unused/PCI space as non-cacheable. More restrictive regions are intended to overlay the 640k to 1MByte area. Failure to program ARR7 with the correct amount of main memory can result in:

- Incorrect memory sizing by the operating system eventually resulting in failure,
- PCI devices not working correctly or causing the system to hang,
- Low performance if ARR7 is programmed with a smaller size than the actual amount of memory.

If the granularity selection in ARR7 does not accommodate the exact size of main memory, unused ARRs can be used to fill-in as non-cacheable regions. All unused/PCI memory space must always be set as non-cacheable.

---

### 5.2.2 *I/O Recovery Time (IORT)*

Back-to-back I/O writes followed by I/O reads may occur too quickly for a peripheral to respond correctly. Historically, programmers have inserted several “JMP \$+2” instructions in the hope that code fetches on the bus would create sufficient recovery time. The 6x86MX’s Branch Target Buffer (BTB) typically eliminates these external code fetches, thus the previous method of guaranteeing I/O recovery no longer applies. For the 6x86MX, one approach to dealing with this issue is to insert I/O write cycles to a dummy port. I/O write cycles in the form of “out imm,reg” are easily implemented as shown below:

OLD IORT	NEW IORT
out 21h,al	out 21h,al
jmp \$+2	out 80h,al
jmp \$+2	out 80h,al
jmp \$+2	out 80h,al
in al,21h	in al,21h

The 6x86MX incorporates an alternative method for implementing I/O recovery time using user selectable delay settings. See the section on 6x86MX IORT settings below.

---

### 5.2.3 *BIOS Creation Utilities*

BIOS creation utilities or setup screens must have the capability to easily define and modify the contents of the 6x86MX configuration registers. This allows OEMs and integrators to easily configure these register settings with the values appropriate for their system design.

---

## 5.3 *Recommended Bit Settings*

---

### 5.3.1 *NC1*

The NC1 bit in CCR0 controls the predefined non-cacheable region from 640K to 1 MByte. The 640K to 1MByte region should be non-cacheable to prevent L1 caching of expansion cards using memory mapped I/O (MMIO). Setting NC1 also implies that the video BIOS and system BIOS are non-cacheable. Experiments with both the 6x86MX and Pentium CPUs have shown that performance is largely unchanged whether the video BIOS and system BIOS was cached or not. This assumes that a modern operating system was used and that the measurements are taken with a recent benchmark applications, such as WinStone95.

Recommended setting: NC1 = 1

---

### 5.3.2 *NO\_LOCK*

NO\_LOCK enables weak locking for the entire address space. NO\_LOCK may cause failures for software that requires locked cycles in order to operate correctly.

Recommended setting: NO\_LOCK = 0

---

### 5.3.3 *LOCK\_NW*

Once set, LOCK\_NW prohibits software from changing the NW bit in CR0. Since the definition of the NW bit is the same for both the 6x86MX and the Pentium, it is not necessary to set this bit.

Recommended setting: LOCK\_NW = 0

---

### 5.3.4 *WPR1*

WPR1 forces cacheable accesses in the 640k to 1MByte address region to be write-protected. If NC1 is set (recommended setting), all caching is disabled from 640k to 1MByte and WPR1 is not required. However, if ROM areas within the 640k-1MByte address region are cached, WPR1 should be set to protect against errant self-modifying code.

Recommended setting: WPR1 = 0 unless ROM areas are cached

---

### 5.3.5 *LINBRST*

Linear Burst (LINBRST) allows for an alternate address sequence for burst cycles. The system logic, L2 cache and motherboard design must also support this feature in order for the 6x86MX to function properly with this bit enabled. Linear Burst provides higher performance than the default “1+4” burst sequence, but should only be enabled if the system is designed to support it.

If the system does support linear burst, BIOS should enable this feature in both the system logic and the 6x86MX prior to enabling the L1 cache. Appendix G includes sample code that can be used to detect if the L2 cache supports linear burst mode.

Recommended setting: LINBRST = 0 unless linear burst supported by the system

---

### 5.3.6 *MAPEN*

When set to 1h, the MAPEN bits allow access to all 6x86MX configuration registers including indices outside the C0h-CFh and FCh-FFh ranges. MAPEN should be set to 1h only to access specific configuration registers and then should be cleared immediately after the access is complete.

Recommended setting: MAPEN(3-0) = 0 except for specific configuration register accesses

---

### 5.3.7 *IORT*

I/O recovery time specifies the minimum number of bus clocks between I/O accesses for the CPU's bus controller. The system logic typically has a built-in method to select the amount of I/O recovery time. It is preferred to configure the system logic with the I/O recovery time setting and set the CPU for a minimum I/O recovery time delay.

Recommended setting: IORT(2-0) = 7

---

### 5.3.8 *CPUID*

When set, the CPUID bit enables the CPUID instruction. By default, the CPUID instruction is enabled (CPUID = 1).

When enabled, the CPUID opcode is enabled and the CPUID bit in the EFLAGS can be modified. The CPUID instruction can then be called to inspect the type of CPU present.

When the CPUID instruction is disabled (CPUID = 0), the CPUID opcode 0FA2 causes an invalid opcode exception. Additionally, the CPUID bit in the EFLAGS register cannot be modified by software.

Recommended setting: CPUID = 1

---

### 5.3.9 *WT\_ALLOC*

Write Allocate (WT\_ALLOC) allows L1 cache write misses to cause a cache line allocation. This feature improves the L1 cache hit rate resulting in higher performance. Especially useful for Windows applications.

Recommended setting: WT\_ALLOC = 1

---

### 5.3.10 *ARREN*

The ARREN bit enables or disables all eight ARR's. When ARREN is cleared (default), the ARR's can be safely programmed. Most systems will need to use at least one address region register (ARR). Therefore, ARREN should always be set after the ARR's and RCR's have been initialized.

Recommended setting: ARREN = 1 after initializing ARR0-ARR7, RCR0-RCR7

---

**Recommended Bit Settings**

---

**5.3.11 ARR7 and RCR7**

Address Region 7 (ARR7) defines the Main Memory Region (MMR). This region specifies the amount of cacheable main memory and its attributes. Once BIOS completes memory discovery, ARR7 should be programmed with a base address of 000 0000h and with a "Size" equal to the amount of main memory installed in the system. Memory accesses outside of this region are defined as non-cacheable to ensure compatibility with PCI devices.

Recommended settings:

ARR7 Base Addr= 0000 0000h

ARR7 Block Size= amount of main memory

RCR7 RCE = 1

RCR7 WL = 0

RCR7 WG = 1

RCR7 WT = 0

If the granularity selection in ARR7 does not accommodate the exact size of main memory, unused ARR7s can be used to fill-in as non-cacheable regions (RCD = 1) as shown in the table below. All unused/PCI memory space must always be set as non-cacheable.

MEM	ARR7		ARR6		ARR5		ARR4	
	BASE (HEX)	SIZE (MB)						
8	0	8						
16	0	16						
24	0	32	0180 0000	8				
32	0	32						
40	0	64	0300 0000	16	0280 0000	8		
48	0	64	0300 0000	16				
64	0	64						
72	0	128	0600 0000	32	0500 0000	16	0480 0000	8
80	0	128	0600 0000	32	0500 0000	16		
96	0	128	0600 0000	32				
128	0	128						
160	0	256	0E00 0000	32	0C00 0000	32	0A00 0000	32
192	0	256	0E00 0000	32	0C00 0000	32		
256	0	256						

**ARR Settings for Various Main Memory Sizes**

---

### 5.3.12 SMM Features

The 6x86MX supports SMM mode through the use of the SMI# and SMIACT# pins, and a dedicated memory region for the SMM address space. SMM features must be enabled prior to servicing any SMI interrupts. The following paragraphs describe each of the SMM features and recommended settings.

#### USE\_SMI

Prior to servicing SMI interrupts, SMM-capable systems must enable the SMM pins by setting USE\_SMI=1. The SMM hardware pins (SMI# and SMIACT#) are disabled by default.

#### SMAC

If set, any access to addresses within the SMM address space are directed to SMM memory instead of main memory. Setting SMAC allows access to the SMM memory without servicing an SMI. Additionally, SMAC allows use of the SMINT instruction (software SMI). This bit may be enabled to initialize or test SMM memory but should be cleared for normal operation.

#### SM3 and ARR3

Address Region Register 3 (ARR3) can be used to define the System Management Address Region (SMAR). Systems that use SMM features must use ARR3 to establish a base and limit for the SMM address space.

Only ARR3 can be used to establish the SMM region.

Typically, SMAR overlaps normal address space. RCR3 defines the attributes for both the SMM address region *and* the normal address space. If SMAR overlaps main memory, write gathering should be enabled for ARR3. If SMAR overlaps video memory, ARR3 should be set as non-cacheable and write gathering should be enabled.

#### NMI\_EN

The NMI\_EN bit allows NMI interrupts to occur within an SMI service routine. If this feature is enabled, the SMI service routine must guarantee that the IDT is initialized properly to allow the NMI to be serviced. Most systems do not require this feature.

## SMI\_LOCK

Once the SMM features are initialized in the configuration registers, they can be permanently locked using the SMI\_LOCK bit. Locking the SMM related bits and registers prevents applications from tampering with these settings. Even if SMM is not implemented, setting SMI\_LOCK in combination with SMAC=0 prevents software SMIs from occurring.

Once SMI\_LOCK is set, it can only be cleared by a processor RESET. Consequently, setting SMI\_LOCK makes system/BIOS/SMM debugging difficult. To alleviate this problem, SMI\_LOCK must be implemented as a user selectable "Secure SMI (enable/disable)" feature in CMOS setup. If SMI\_LOCK is not user selectable, it is recommended that SMI\_LOCK = 0 to allow for system debug.

Suggested settings for systems not using SMM:

USE_SMI	= 0
SMAC	= 0
SM3	= 0
ARR3	= may be used as normal address region register
SMI_LOCK	= 0
NMI_EN	= 0

Suggested settings for systems using SMM:

USE_SMI	= 1
SMAC	= 0
SM3	= 1
ARR3 Base Addr	= as required
ARR3 Block Size	= as required
SMI_LOCK	= 0
NMI_EN	= 0

---

### 5.3.13 Power Management Features

#### SUSP\_HALT

Suspend on Halt (SUSP\_HLT) permits the CPU to enter a low power suspend mode when a HLT instruction is executed. Although this provides some power management capability, it is not optimal.

---

## Recommended Bit Settings

Suggested setting:  
SUSP\_HALT = 0

### USE\_SUSP

In addition to the HLT instruction, low power suspend mode may be activated using the SUSP# input pin. In response to the SUSP# input, the SUSPA# output indicates when the 6x86MX has entered low power suspend mode. Systems that support the 6x86MX's low power suspend feature via the hardware pins must set USE\_SUSP to enable these pins.

Suggested setting:  
USE\_SUSP = 0 unless hardware suspend pins supported

---

## 6. Model Specific Registers

The 6x86MX contains four model specific registers (MSR0 - MSR3). These 64-bit registers are listed in the table below.

**MACHINE SPECIFIC REGISTER**

REGISTER DESCRIPTION	MSR ADDRESS	REGISTER
Time Stamp Counter (TSC)	10h	MSR10
Counter Event Selection and Control Register	11h	MSR11
Performance Counter #0	12h	MSR12
Performance Counter #1	13h	MSR13

The MSR registers can be read using the RDMSR instruction, opcode 0F32h. During an MSR register read, the contents of the particular MSR register, specified by the ECX register, is loaded into the EDX:EAX registers.

The MSR registers can be written using the WRMSR instruction, opcode 0F30h. During a MSR register write the contents of EDX:EAX are loaded into the MSR register specified in the ECX register.

The RDMSR and WRMSR instructions are privileged instructions.

---

### 6.1 Time Stamp Counter

The Time Stamp Counter (TSC) Register (MSR10) is a 64-bit counter that counts the internal CPU clock cycles since the last reset. The TSC uses a continuous CPU core clock and will continue to count clock cycles even when the 6x86MX is suspend mode or shutdown.

The TSC can be accessed using the RDMSR and WRMSR instructions. In addition, the TSC can be read using the RDTSC instruction, opcode 0F31h. The RDTSC instruction loads the contents of the TSC into EDX:EAX. The use of the RDTSC

instruction is restricted by the Time Stamp Disable, (TSD) flag in CR4. When the TSD flag is 0, the RDTSC instruction can be executed at any privilege level. When the TSD flag is 1, the RDTSC instruction can only be executed at privilege level 0.

---

## *6.2 Performance Monitoring*

Performance monitoring allows counting of over a hundred different event occurrences and durations. Two 48-bit counters are used: Performance Monitor Counter 0 and Performance Monitor Counter 1. These two performance monitor counters are controlled by the Counter Event Control Register (MSR11). The performance monitor counters use a continuous CPU core clock and will continue to count clock cycles even when the 6x86MX is in suspend mode or shutdown.

---

## *6.3 Performance Monitoring Counters 1 and 2*

The 48-bit Performance Monitoring Counters (PMC) Registers (MSR12, MSR13) count events as specified by the counter event control register. The PMCs can be accessed by the RDMSR and WRMSR instructions. In addition, the PMCs can be read by the RDPMC instruction, opcode 0F33h. The RDPMC instruction loads the contents of the PMC register specified in the ECX register into EDX:EAX. The use of RDPMC instructions is restricted by the Performance Monitoring Counter Enable, (PCE) flag in C4. When the PCE flag is set to 1, the RDPMC instruction can be executed at any privilege level. When the PCE flag is 0, the RDPMC instruction can only be executed at privilege level 0.

---

## *6.4 Counter Event Control Register*

Register MSR 11h controls the two internal counters, #0 and #1. The events to be counted have been chosen based on the micro-architecture of the 6x86MX processor. The control register for the two event counters is described on page 46.

---

## *6.5 PM Pin Control*

The Counter Event Control register (MSR11) contains PM control fields that define the PM0 and PM1 pins as counter overflow indicators or counter event indicators. When defined as event counters, the PM pins indicate that one or more events

occurred during a particular clock cycle and do not count the actual events. When defined as overflow indicators, the event counters can be preset with a value less than  $2^{48}-1$  and allowed to increment as events occur. When the counter overflows the PM pin becomes asserted.

---

### *6.5.1 Counter Type Control*

The Counter Type bit determines whether the counter will count clocks or events. When counting clocks the counter operates as a timer.

---

### *6.5.2 CPL Control*

The Current Privilege Level (CPL) can be used to determine if the counters are enabled. The CP02 bit in the MSR 11 register enables counting when the CPL is less than three, and the CP03 bit enables counting when CPL is equal to three. If both bits are set, counting is not dependent on the CPL level; if neither bit is set, counting is disabled.



### 6.5.3 Event Type and Description

The events that can be counted by the performance monitoring counters are listed in Figure 1-. Each of the 127 event types is assigned an event number. A particular event number to be counted is placed in one of the MSR 11 Event Type fields. There is a separate field for counter #0 and #1.

The events are divided into two groups. The occurrence type events and duration type events. The occurrence type events, such as hardware interrupts, are counted as single events. The duration type events such as “clock while bus cycles are in progress” count the number of clock cycles that occur during the event.

During occurrence type events, the PM pins are configured to indicate the counter has incremented. The PM pins will then assert every time the counter increments in regards to an occurrence event. Under the same PM control, for a duration event the PM pin will stay asserted for the duration of the event.

**EVENT TYPE REGISTER**

NUMBER	COUNTER 0	COUNTER 1	DESCRIPTION	TYPE
00h	yes	yes	Data Reads	Occurrence
01h	yes	yes	Data Writes	Occurrence
02h	yes	yes	Data TLB Misses	Occurrence
03h	yes	yes	Cache Misses: Data Reads	Occurrence
04h	yes	yes	Cache Misses: Data Writes	Occurrence
05h	yes	yes	Data Writes that hit on Modified or Exclusive Liens	Occurrence
06h	yes	yes	Data Cache Lines Written Back	Occurrence
07h	yes	yes	External Inquiries	Occurrence
08h	yes	yes	External Inquires that hit	Occurrence
09h	yes	yes	Memory Accesses in both pipes	Occurrence
0Ah	yes	yes	Cache Bank conflicts	Occurrence
0Bh	yes	yes	Misaligned data references	Occurrence
0Ch	yes	yes	Instruction Fetch Requests	Occurrence
0Dh	yes	yes	L2 TLB Code Misses	Occurrence
0Eh	yes	yes	Cache Misses: Instruction Fetch	Occurrence
0Fh	yes	yes	Any Segment Register Load	Occurrence
10h	yes	yes	Reserved	Occurrence
11h	yes	yes	Reserved	Occurrence
12h	yes	yes	Any Branch	Occurrence
13h	yes	yes	BTB hits	Occurrence
14h	yes	yes	Taken Branches or BTB hits	Occurrence
15h	yes	yes	Pipeline Flushes	Occurrence
16h	yes	yes	Instructions executed in both pipes	Occurrence
17h	yes	yes	Instructions executed in Y pipe	Occurrence
18h	yes	yes	Clocks while bus cycles are in progress	Duration
19h	yes	yes	Pipe Stalled by full write buffers	Duration
1Ah	yes	yes	Pipe Stalled by waiting on data memory reads	Duration

## EVENT TYPE REGISTER (CONTINUED)

NUMBER	COUNTER 0	COUNTER 1	DESCRIPTION	TYPE
1Bh	yes	yes	Pipe Stalled by writes to not-Modified or not-Exclusive cache lines.	Duration
1Ch	yes	yes	Locked Bus Cycles	Occurrence
1Dh	yes	yes	I/O Cycles	Occurrence
1Eh	yes	yes	Non-cacheable Memory Requests	Occurrence
1Fh	yes	yes	Pipe Stalled by Address Generation Interlock	Duration
20h	yes	yes	Reserved	
21h	yes	yes	Reserved	
22h	yes	yes	Floating Point Operations	Occurrence
23h	yes	yes	Breakpoint Matches on DR0 register	Occurrence
24h	yes	yes	Breakpoint Matches on DR1 register	Occurrence
25h	yes	yes	Breakpoint Matches on DR2 register	Occurrence
26h	yes	yes	Breakpoint Matches on DR3 register	Occurrence
27h	yes	yes	Hardware Interrupts	Occurrence
28h	yes	yes	Data Reads or Data Writes	Occurrence
29h	yes	yes	Data Read Misses or Data Write Misses	Occurrence
2Bh	yes	no	MMX Instruction Executed in X pipe	Occurrence
2Bh	no	yes	MMX Instruction Executed in Y pipe	Occurrence
2Dh	yes	no	EMMS Instruction Executed	Occurrence
2Dh	no	yes	Transition Between MMX Instruction and FP Instructions	Occurrence
2Eh	no	yes	Reserved	
2Fh	yes	no	Saturating MMX Instructions Executed	Occurrence
2Fh	no	yes	Saturations Performed	Occurrence
30h	yes	no	Reserved	
31h	yes	no	MMX Instruction Data Reads	Occurrence
32h	yes	no	Reserved	
32h	no	yes	Taken Branches	Occurrence
33h	no	yes	Reserved	
34h	yes	no	Reserved	
34h	no	yes	Reserved	
35h	yes	no	Reserved	
35h	no	yes	Reserved	
36	yes	no	Reserved	
36	no	yes	Reserved	
37	yes	no	Returns Predicted Incorrectly	Occurrence
37	no	yes	Return Predicted (Correctly and Incorrectly)	Occurrence
38	yes	no	MMX Instruction Multiply Unit Interlock	Duration
38	no	yes	MODV/MOVQ Store Stall Due to Previous Operation	Duration
39	yes	no	Returns	Occurrence
39	no	yes	RSB Overflows	Occurrence
3A	yes	no	BTB False Entries	Occurrence

**EVENT TYPE REGISTER (CONTINUED)**

<b>NUMBER</b>	<b>COUNTER 0</b>	<b>COUNTER 1</b>	<b>DESCRIPTION</b>	<b>TYPE</b>
3A	no	yes	BTB Miss Prediction on a Not-Taken Back	Occurrence
3B	yes	no	Number of Clock Stalled Due to Full Write Buffers While Executing	Duration
3B	no	yes	Stall on MMX Instruction Write to E or M Line	Duration
3C - 3Fh	yes	yes	Reserved	Duration
40h	yes	yes	L2 TLB Misses (Code or Data)	Occurrence
41h	yes	yes	L1 TLB Data Miss	Occurrence
42h	yes	yes	L1 TLB Code Miss	Occurrence
43h	yes	yes	L1 TLB Miss (Code or Data)	Occurrence
44h	yes	yes	TLB Flushes	Occurrence
45h	yes	yes	TLB Page Invalidates	Occurrence
46h	yes	yes	TLB Page Invalidates that hit	Occurrence
47h	yes	yes	Reserved	
48h	yes	yes	Instructions Decoded	Occurrence
49h	yes	yes	Reserved	

---

## *7. Programming Model Differences*

---

### *7.1 Instruction Set*

The 6x86MX supports the Pentium Pro instruction set plus MMX instructions. Pentium extensions for virtual mode are not supported.

---

### *7.2 Configuring Internal 6x86MX Features*

The 6x86MX supports configuring internal features through I/O ports.

---

### *7.3 INVD and WBINVD Instructions*

The INVD and WBINVD instructions are used to invalidate the contents of the internal and external caches. The WBINVD instruction first writes back any modified lines in the cache and then invalidates the contents. It ensures that cache coherency with system memory is maintained regardless of the cache operating mode. Following invalidation of the internal cache, the CPU generates special bus cycles to indicate that external caches should also write back modified data and invalidate their contents.

On the 6x86MX, the INVD functions identically to the WBINVD instruction. The 6x86MX always writes all modified internal cache data to external memory prior to invalidating the internal cache contents. In contrast, the Pentium invalidates the contents of its internal caches without writing back the “dirty” data to system memory.

---

#### *7.4 Control Register 0 (CR0) CD and NW Bits*

The CPU's CR0 register contains, among other things, the CD and NW bits which are used to control the on-chip cache. CR0, like the other system level registers, is only accessible to programs running at the highest privilege level. The table on the following page lists the cache operating modes for all possible states of the CD and NW bits.

The CD and NW bits are set to one (cache disabled) after reset. For highest performance the cache should be enabled in write-back mode by clearing the CD and NW bits to 0. Sample code for enabling the cache is listed in Appendix E. To completely disable the cache, it is recommended that CD and NW be set to 1 followed by execution of the WBINVD instruction. The 6x86MX cache always accepts invalidation cycles even when the cache is disabled. Setting CD=0 and NW=1 causes a General Protection fault on the Pentium, but is allowed on the 6x86MX to globally enable write-through caching.

---

**Control Register 0 (CR0) CD and NW Bits**

<b>CD</b>	<b>NW</b>	<b>OPERATING MODES</b>
1	1	Cache disabled. Read hits access the cache. Read misses do not cause line fills. Write hits update the cache and system memory. Write hits change exclusive lines to modified. Shared lines remain shared after write hit. Write misses access memory. Inquiry and invalidation cycles are allowed. System memory coherency maintained.
1	0	Cache disabled. Read hits access the cache. Read misses do not cause line fills. Write hits update the cache. Only write hits to shared lines and write misses update system memory. Write misses access memory. Inquiry and invalidation cycles are allowed. System memory coherency maintained.
0	1	Cache enabled in Write-through mode. Read hits access the cache. Read misses may cause line fills. Write hits update the cache and system memory. Write misses access memory. Inquiry and invalidation cycles are allowed. System memory coherency maintained.
0	0	Cache enabled in Write-back mode. Read hits access the cache. Read misses may cause line fills. Write hits update the cache. Write misses access memory and may cause line fills if write allocation is enabled. Inquiry and invalidation cycles are allowed. System memory coherency maintained.

**Cache Operating Modes**

---

## *Appendixes*

---

### *Appendix A.- Sample Code: Detecting a Cyrix CPU*

```
assume cs:_TEXT
public _iscyrix
_TEXT segment byte public 'CODE'
;*****
;      Function: int iscyrix ()
;
;      Purpose:      Determine if Cyrix CPU is present
;      Technique:    Cyrix CPUs do not change flags where flags
;                   change in an undefined manner on other CPUs
;
;      Inputs:       none
;      Output:       ax == 1 Cyrix present, 0 if not
;*****
_iscyrixproc near
    .386
    xor  ax, ax          ; clear ax
    sahf                 ; clear flags, bit 1 always=1 in flags
    mov  ax, 5
    mov  bx, 2
    div  bl              ; operation that doesn't change flags
    lahf                 ; get flags
    cmp  ah, 2           ; check for change in flags
    jne  not_cyrix      ; flags changed, therefore NOT CYRIX
    mov  ax, 1          ; TRUE Cyrix CPU
    jmp  done

not_cyrix:
    mov  ax, 0          ; FALSE NON-Cyrix CPU
done:
    ret
_iscyrix endp
_TEXT ends
end
```

---

**Sample Code: Determining CPU MHz**

---

*Appendix B. Sample Code: Determining CPU MHz*

```
assume cs:_TEXT
public _cpu_speed
_TEXT segment para public 'CODE'

comment~
*****
Function:      unsigned long _cpu_speed( unsigned int )
               "C" style caller
Purpose:       calculate elapsed time req'd to complete a loop of IDIVs

Technique:    Use the PC's high resolution timer/counter chip (8254)
               to measure elapsed time of a software loop consisting
               of the IDIV and LOOP instruction.

Definitions:  The 8254 receives a 1.19318MHz clock (0.8380966 usec).
               One "tick" is equal to one rising clock edge applied
               to the 8254 clock input.

Inputs:       ax = no. of loops for cpu_speed_loop
Returns:      ax = no. of 1.19318MHz clk ticks req'd to complete a loop
               dx = state of 8254 out pin
*****~
PortB         EQU      061h
Timer_Ctrl_Reg EQU      043h
Timer_2_Data  EQU      042h
stk$dx        EQU      10          ;dx register offset
stk$ax        EQU      14          ;dx register offset
stack$ax      EQU      [bp]+stk$ax
stack$dx      EQU      [bp]+stk$dx
Loop_Count    EQU      [bp+16]+4

.386p

_cpu_speed    proc near
    pushf                    ;save interrupt flag
    pusha                   ;pushes 16 bytes on stack
    mov     bp,sp           ;init base ptr

    cli                      ;disable interrupts

;-----disable clock to timer/counter 2
    in     al, PortB
    and   al, 0feh
    out   80h,al            ;I/O recovery time
    out   PortB, al
    mov   di, ax
```

```
;-----initialize the 8254 counter to "0", known value
    mov     al,0b0h
    out     Timer_Ctrl_Reg, al      ;control word to set channel 2 count
    out     80h,al                 ;I/O recovery time
    mov     al,0ffh
    out     Timer_2_Data, al       ;init count to 0, lsb
    out     80h,al                 ;I/O recovery time
    out     Timer_2_Data, al       ;init count to 0, msb

;-----get the number of loops from the caller's stack
    mov     cx,Loop_Count          ;loop count

;-----load dividend & divisor, clk count for IDIV depend on operands!
    xor     edx,edx                ;dividend EDX:EAX
    xor     eax,eax
    mov     ebx,1                  ;divisor

;-----enable the timer/counter's clock. Begin timed portion of test!
    xchg   ax, di                  ;save ax for moment
    or     al, 1
    out    PortB, al               ;enable timer/counter 2 clk
    xchg   ax, di                  ;restore ax

;-----this is the core loop.
    ALIGN 16
cpu_speed_loop:
    idiv   ebx
    idiv   ebx
    idiv   ebx
    idiv   ebx
    idiv   ebx
    loop   cpu_speed_loop

;-----disable the timer/counter's clk. End timed portion of test!
    mov    ax, di
    and    al, 0FEH
    out    PortB, al

;-----send latch status command to the timer/counter
    mov    al, 0c8h                ;latch status and count
    out    Timer_Ctrl_Reg, al
    out    80h,al                 ;I/O recovery time

;-----read status byte, and count value "ticks" from the timer/cntr
    in     al, Timer_2_Data         ;read status
    out    80h,al                 ;I/O recovery time
    mov    dl, al
    and    dx, 080h
    shr    dx, 7
```

---

**Sample Code: Determining CPU MHz**

```

    in     al, Timer_2_Data      ;read LSB
    out   80h,al                ;I/O recovery time
    mov   bl, al
    in   al, Timer_2_Data      ;read MSB
    out   80h,al                ;I/O recovery time
    mov   bh, al

    not   bx                    ;invert count

;-----send command to clear the timer/counter
    mov   al, 0b6h
    out   Timer_Ctrl_Reg, al    ;clear channel 2 count
    out   80h,al                ;I/O recovery time
    xor   al, al
    out   Timer_2_Data, al      ;set count to 0, lsb
    out   80h,al                ;I/O recovery time
    out   Timer_2_Data, al      ;set count to 0, msb

;-----put return values on the stack for the caller
    mov   [bp+stk$ax], bx
    mov   [bp+stk$dx], dx

    popa
    popf                        ;restores interrupt flag
    ret

_cpu_speed    endp

.8086
_TEXT    ENDS
END
```

---

*Appendix C. Example CPU Type and Frequency Detection Program*

```

/* *****
function:      main()                                     WCP 8/22/95
Purpose:      a driver program to demonstrate:
              CPU detection
              CPU core frequency in Mhz.
Returns:      0 if successful

Required source code modules
ml_stat.c      main() module (this file)
id.asm         cpu identification code
clock.asm      cpu timing loop

Compile and Link instructions for Borland C++ or equivalent:
    bcc ml_stat.c id.asm clock.asm
***** */
/* include directives */
#include <stdio.h>

/* constants */
#define TTPS      1193182 //high speed Timer Ticks per second in Mhz
#define MHZ      1000000 //number of clocks in 1 Mhz
#define LOOP_COUNT 0x2000 //core loop iterations
#define RUNS     10 //number of runs to average
#define DIVS     5 //# of IDIV instructions in the core loop
#define 6x86MX_IDIV_CLKS 17 //known clock counts for 6x86MX
#define 6x86MX_LOOP_CLKS 1
#define P54_IDIV_CLKS 46 //known clock counts for P54
#define P54_LOOP_CLKS 7

/* prototypes */
unsigned int  iscyrix( void ); //detects cyrix cpu
unsigned long cpu_speed( unsigned int ); //core timing loop

main(){

/* declarations */
unsigned char uc_cyrix_cpu = 0; //Cyrix cpu? 0=no, 1=yes
unsigned int i_runs = 0; //number of runs to avg
unsigned int ui_idiv, ui_loop = 0; //instruction clk counts
unsigned long ul_tt_cnt, ul_tt_sum = 0; //timer tick counts, sum
unsigned int ui_core_loop_cntr = LOOP_COUNT; //core loop iterations
float f_mtt = 0; //measured timer ticks
float f_total_core_clks = 0; //calculated core clocks
float f_total_time = 0; //measured time
float f_mhz = 0; //mhz

```

---

**Example CPU Type and Frequency Detection Program**

```
/* ***** determine if Cyrix CPU is present ***** */
//detect if Cyrix CPU is present
uc_cyrix_cpu = iscyrix(); //1=cyrix, 0=non-cyrix

//display a msg
if(uc_cyrix_cpu) printf("\nCyrix CPU present! ");
else printf("\nCyrix CPU not present! ");

/* ***** determine CPU Mhz ***** */

//count # of hi speed "timer ticks" to complete several runs of core loop
for (i_runs = 0 ; i_runs < RUNS ; i_runs++) {
    ul_tt_cnt = cpu_speed( ui_core_loop_cntr );
    ul_tt_sum += ul_tt_cnt; //sum them all together
} //end for

//compute the avg number of high speed "timer ticks" for the several runs
f_mtt = ul_tt_sum / RUNS; //compute the average

//initialize variables with the "known" clock counts for a 6x86MX or P54
if(uc_cyrix_cpu)ui_idiv=6x86MX_IDIV_CLKS; else ui_idiv=P54_IDIV_CLKS;
if(uc_cyrix_cpu)ui_loop=6x86MX_LOOP_CLKS; else ui_loop=P54_LOOP_CLKS;

//determine the total number of core clocks. (5 idivs are in the core
loop)
f_total_core_clks = (float)ui_core_loop_cntr * (ui_idiv * DIVS + ui_loop);

//the time it took to complete the core loop can be determined by the
//ratio of measured timer ticks(mtt) to timer ticks per second(TTPS).
f_total_time = f_mtt / TTPS;

//frequency can be found by the ratio of core clks to the total time.
f_mhz = f_total_core_clks / f_total_time;
f_mhz = f_mhz / MHZ; //convert to Mhz

//display a msg
printf("The core clock frequency is: %3.1f MHz\n\n",f_mhz);

return(0);
} //end main
```

---

## *Appendix D.- Sample Code: Programming 6x86MX Configuration Registers*

### *Reading/Writing Configuration Registers*

Sample code for setting NC1=1 in CCR0.

```
pushf                ;save the if flag
cli                  ;disable interrupts
mov  al, 0c0h        ;set index for CCR0
out  22h, al         ;select CCR0 register
in   al, 23h         ;READ current CCR0 valueREAD

mov  ah, al
or   ah, 2h          ;MODIFY, set NC1 bitMODIFY

mov  al, 0c0h        ;set index for CCR0
out  22h, al         ;select CCR0 register
mov  al, ah
out  23h,al          ;WRITE new value to CCR0WRITE
popf                  ;restore if flag
```

### *Setting MAPEN*

Sample code for setting MAPEN=1 in CCR3 to allow access to all the configuration registers.

```
pushf                ;save the if flag
cli                  ;disable interrupts
mov  al, 0c3h        ;set index for CCR3
out  22h, al         ;select CCR3 register
in   al, 23h         ;current CCR3 valueREAD

mov  ah, al
and  ah,0Fh          ;clear upper nibble of ah
or   ah, 10h         ;MODIFY, set MAPEN(3-0)MODIFY

mov  al, 0c3h        ;set index for CCR3
out  22h, al         ;select CCR3 register
mov  al, ah
out  23h,al          ;WRITE new value to CCR3WRITE
popf                  ;restore if flag
```

---

**- Sample Code: Controlling the L1 Cache**

---

*Appendix E. - Sample Code: Controlling the L1 Cache*

*Enabling the L1 Cache*

;reading/writing CR0 is a privileged operation.

```
mov    eax, cr0
and    eax, 09fffffffh;clear the CD=0, NW=1 bits to enable write-back
mov    cr0, eax;control register 0 write
wbinvd                ;optional, by flushing the L1 cache here it
                        ;ensures the L1 cache is completely clean
```

*Disabling the L1 Cache*

```
mov    eax, cr0
or     eax, 060000000h ;set the CD=1, NW=1 bits to disable caching
mov    cr0, eax      ;control register 0 write
wbinvd
```

*Appendix F. - Example Configuration Register Settings*

Below is an example of optimized 6x86MX settings for a 16 MByte system with PCI. Since SMI address space overlaps Video RAM at A0000h, WG is set to maintain the settings of the underlying region ARR0. If SMI address space overlapped system memory at 30000h, only WWO and WG would be set. If SMI address space overlapped FLASH ROM at E0000h, only RCD would be set. Power management features are disabled in this example system.

**CONFIGURATION REGISTER SETTINGS EXAMPLE**

REGISTER	BIT(S)	SETTING	DESCRIPTION
CCR0	NC1	1	Disables caching from 640k-1MByte.
CCR1	USE_SMI	1	Enables SMI# and SMIACT# pins.
	SMAC	0	Always clear SMAC for normal operation.
	NO_LOCK	0	Enforces strong locking for compatibility.
	SM3	1	Sets ARR3 as SMM address region.
CCR2	LOCK_NW	0	Locking NW bit not required.
	SUSP_HLT	0	Power management not required for this system.
	WPR1	0	ROM areas not cached, so WPR1 not required.
	USE_SUSP	0	Power management not required for this system.
CCR3	SMI_LOCK	0	Locks SMI feature as initialized.
	NMI_EN	0	Servicing NMIs during SMI not required.
	LINBRST	0	Linear burst not supported in this system.
	MAPEN(3-0)	0	Always clear MAPEN for normal operation.
CCR4	IORT(2-0)	7	Sets IORT to minimum setting.
	CPUIDEN	1	Enables CPUID instruction.
CCR5	WT_ALLOC	1	Enables write allocation for performance.
	ARREN	1	Enables all ARR's.
ARR0	BASE ADDR	A0000h	Video buffer base address = A0000h.
	BLOCK SIZE	6h	Video buffer block size = 128KBytes.
RCR0	RCD	1	Caching disabled for compatibility. Caching also disabled via NC1.
	WL	0	
	WG	1	
	WT	0	Write gathering enabled for performance.
	INV_RGN	0	

---

- Example Configuration Register Settings

**CONFIGURATION REGISTER SETTINGS EXAMPLE (CONTINUED)**

ARR1	BASE ADDR	C 0000h	Expansion Card/ ROM base address = C 0000h.
	BLOCK SIZE	7h	Expansion Card/ROM block size = 256KBytes.
RCR1	RCD	1	Caching disabled for compatibility. Caching also disabled via NC1.
	WL	0	
	WG	0	
	WT	0	
	INV_RGN	0	
ARR3	BASE ADDR	A0000h	SMM address region base address
	BLOCK SIZE	4h	SMM address space = 32 KBytes
RCR3	RCD	1	Caching disabled due to overlap with video buffer.
	WL	0	
	WG	1	Write gathering enabled due to overlap with video buffer.
	WT	0	
	INV_RGN	0	
ARR7	BASE ADDR	0h	Main memory base address = 0h.
	BLOCK SIZE	7h	Main memory size = 16 MBytes.
RCR7	RCE	1	Caching, write gathering enabled for main memory.
	WL	0	
	WG	1	
	WT	0	
ARR(2,4-6)	BASE ADDR	0	ARR(2,4-6) disabled (default state).
	BLOCK SIZE	0	
RCR(2,4-6)	RCD	0	RCR(2,4-6) not required due to corresponding ARR's disabled (default state).
	WL	0	
	WG	0	
	WT	0	
	INV_RGN	0	

---

*Appendix G. - Sample Code: Detecting L2 Cache Burst Mode*

comment~\*\*\*\*\*

Purpose: This example program detects if Linear Burst mode is supported.

Method: There are 3 components (CPU, chipset, SPBSRAM) that must agree on the burst order. The CPU and chipset burst order can be determined by inspecting each devices internal configuration registers. The SPBSRAM devices must be interrogated by a software algorithm (below) to determine if "linear burst mode" is enabled/supported correctly.

Algorithm: If the CPU and chipset are programmed for linear burst mode and a known data pattern exists in memory, then the burst mode of the SPBSRAMs can be determined by performing a cache line burst and then inspect the data pattern.

Application: In this example, the SIS5511 chipset is used with a Cyrix 6x86MX CPU.

Environment: This program is a REAL mode DOS program to serve as an example. This example algorithm should be ported to BIOS.

Warnings: For simplicity, this program does not check to see which CPU or chipset is present. Nor, does this program check to see if the CPU is in REAL mode before executing protected instructions. Also, this program blindly overwrites data in the 8000h segment of memory.

\*\*\*\*\*~

```
;version m510 ;remove comment for TASM

DOSSEG
.MODEL SMALL
.DATA
Msg_1 db 0dh,0ah
db 'ISLINBUR.EXE checks if L2 SRAMs are in Linear Burst Mode
or '
db 0dh,0ah
db 'Toggle Burst mode for the SIS5511 chipset and the 6x86MX
CPU.'
db 0dh,0ah
db '$'
```

---

**- Sample Code: Detecting L2 Cache Burst Mode**

```
Msg_2      db      0dh,0ah
           db      'Test complete!'
           db      0dh,0ah
           db      '$'
Msg_yes    db      0dh,0ah
           db      'The L2 SRAMs correctly operate in linear burst mode.'
           db      0dh,0ah
           db      '$'
Msg_no     db      0dh,0ah
           db      'ERROR: The L2 SRAMs incorrectly operate in linear burst
mode.'
           db      0dh,0ah
           db      '$'

index_port dw      0CF8h
data_port  dw      0CFCh
pci_index  dd      80000000h

.STACK 100h
.CODE
.STARTUP
.486P

        pushf
        cli

;-----display a msg using a DOS call
        mov     ax,seg Msg_1
        mov     ds,ax
        mov     dx,offset Msg_1      ;set msg_1 start
        mov     ah,9h                ;print string function
        int     21h                  ;DOS int

;-----disable the L1 internal cache
        call    cache_off
        out     80h,a1                ;write to PC diagnostic port

;-----setup a work space in main memory to perform burst
;mode tests and initialize the memory work space with a
;known pattern
        push    ds
```

```
mov     ax,8000h           ;choose segment 8000h
mov     ds,ax
mov     al,0001h
mov     byte ptr ds:[0],al ;init memory locations
inc     al
mov     byte ptr ds:[8h],al
inc     al
mov     byte ptr ds:[10h],al
inc     al
mov     byte ptr ds:[18h],al
pop     ds

;-----enable the SiS5511 chipset's linear burst mode
mov     al,51h             ;al=reg to read
call    r_pci_reg         ;READ al=reg contents
mov     ah,al
or      ah,8              ;MODIFY set linbrst bit
mov     al,51h
call    w_pci_reg         ;WRITE

;-----enable the CPU's linear burst mode
call    en_linbrst

;-----enable L1 caching
call    cache_on

;-----burst several cache lines so that address 80000h is
;in the L2 cache, but NOT in the L1 cache.
push    ds
mov     ax,8000h           ;choose segment 8000h
mov     ds,ax
mov     al,byte ptr ds:[0h] ;line fill to L2 and L1
mov     al,byte ptr ds:[1000h] ;fill L1 line 1
mov     al,byte ptr ds:[2000h] ;fill L1 line 1
mov     al,byte ptr ds:[3000h] ;fill L1 line 1
mov     al,byte ptr ds:[4000h] ;fill L1 line 1,
;now 80000h exists only in the
;L2 cache (not in L1 anymore!)

;-----burst a cache line so that address 80000h will hit
; the L2 cache SRAMs
```

---

**- Sample Code: Detecting L2 Cache Burst Mode**

```
mov     al,byte ptr ds:[8h]
        ;***** Burst Pattern Table *****
        ;if SRAMs in linear burst mode, then
        ;L1 will be filled with:
        ; byte data
        ; 0    01h
        ; 8    02h
        ; 10   03h
        ; 18   04h
        ;if SRAMs in toggle burst mode, then
        ;L1 will be filled with:
        ; byte data
        ; 0    03h
        ; 8    02h
        ; 10   01h
        ; 18   04h

;-----Compare the cache line to the Burst Pattern Table
;above. The signature of the pattern will determine
;if the burst was linear or toggle.

        ;check byte ds:[10] in the L1
mov     al, byte ptr ds:[10h]
        ;it will be a 1 if toggle mode

cmp     al,3h
        ;it will be a 3 if linear mode

pop     ds
jnz     not_linear

is_linear:
mov     dx,offset Msg_yes ;SRAMs in linear burst mode
jmp     over_not
not_linear:
mov     dx,offset Msg_no  ;SRAMs in toggle burst mode
over_not:
wbinvd

;-----disable L1 internal cache
call    cache_off

;-----restore chipset to toggle mode burst order
```

```
    mov     al,51h                ;al=reg to read
    call    r_pci_reg             ;READ al=reg contents
    mov     ah,al
    and     ah,0f7h              ;MODIFY clr linbrst bit
    mov     al,51h
    call    w_pci_reg             ;WRITE

    call    dis_linbrst

;-----restore L1 caching
    call    cache_on

done:
    popf

;-----display a msg using a DOS call
    mov     ax,seg Msg_2
    mov     ds,ax
    mov     ah,9h                ;print string function
    int     21h                  ;DOS int

;-----return to the operating system
.EXIT

comment~*****~
function    r_pci_reg
purpose    read the pci register at the index in al
inputs     al= the index of the pci register
returns    al= the data read from the pci reg
*****~
r_pci_reg PROC

    pushf
    push    eax
    push    dx
    cli

    mov     dx,index_port
    and     eax,0FFh
    or      eax,pci_index
```

---

**- Sample Code: Detecting L2 Cache Burst Mode**

```
    out    dx, eax

    and    al, 3
    mov    dx, data_port
    add    dl, al
    in     al, dx
    xchg   al, bl                ;preserve rtn value

    mov    eax, pci_index
    mov    dx, index_port
    out    dx, eax

    pop    dx
    pop    eax
    popf

    xchg   al, bl
    ret
```

r\_pci\_reg ENDP

```
comment~*****
function    w_pci_reg
inputs      al= the index of the pci register
            ah= the data to write
outputs     modifies chipset registers directly
returns     none
*****~
w_pci_reg proc
```

```
    pushf
    push   eax
    push   bx
    push   dx
    cli

    mov    bx, ax                ;preserve input value(s)

    mov    dx, index_port
    and    eax, 0FFh
```

```
    or      eax,pci_index
    out     dx,eax

    and     al,3
    mov     dx,data_port
    add     dl,al
    mov     al,bh          ;recall data to write
    out     dx,al

    mov     eax,pci_index
    mov     dx,index_port
    out     dx,eax

    pop     dx
    pop     bx
    pop     eax
    popf
    ret
w_pci_reg ENDP

comment~*****
function      en_linbrst
purpose       enable the 6x86MX linbrst bit
inputs        none
outputs       modifies the 6x86MX CPU registers directly
returns       none
*****~
en_linbrst PROC

    mov     ax,0C3C3h          ;set LINBRST
    out     22h,al
    in      al,23h
    xchg    ah,al
    or      ah,4
    out     22h,al
    xchg    ah,al
    out     23h,al

    ret
en_linbrst ENDP
```

---

**- Sample Code: Detecting L2 Cache Burst Mode**

```
comment~*****~
function      dis_linbrst
purpose      disable the 6x86MX linbrst bit
inputs       none
outputs      modifies the 6x86MX CPU registers directly
returns      none
*****~
```

```
dis_linbrst PROC
```

```
    mov     ax,0C3C3h
    out     22h,al
    in      al,23h
    xchg    ah,al
    and     ah,0fbh      ;clear the linbrst bit
    out     22h,al
    xchg    ah,al
    out     23h,al
```

```
    ret
```

```
dis_linbrst ENDP
```

```
comment~*****~
function      cache_off
purpose      disables the L1 cache
inputs       none
returns      none
*****~
```

```
cache_off PROC
```

```
    pushf
    push    eax
    cli
    mov     eax,cr0
    or      eax,60000000h
    mov     cr0,eax
    wbinvd
    jmp     $+2pop     eax
    popf
    ret
```

```
cache_off ENDP
```

```
comment~*****~  
function      cache_on  
purpose       enables the L1 cache  
inputs        none  
returns       none  
*****~
```

```
cache_on PROC  
    pushf  
    push    eax  
    cli  
    mov     eax,cr0  
    and     eax,9FFFFFFFh  
    mov     cr0,eax  
    pop     eax  
    popf  
    ret
```

```
cache_on ENDP
```

```
END
```

---

©1997 Copyright Cyrix Corporation. All rights reserved.

Printed in the United States of America

Trademark Acknowledgments:

Cyrix is a registered trademark of Cyrix Corporation.

Cx486DX, Cx486DX2, Cx486DX4, 5x86, 6x86 and 6x86MX are trademarks of Cyrix Corporation.

Product names used in this publication are for identification purposes only and may be trademarks of their respective companies.

Order Number: 94xxx-xx

Cyrix Corporation

2703 North Central Expressway

Richardson, Texas 75080-2010

United States of America

Cyrix Corporation (Cyrix) reserves the right to make changes in the devices or specifications described herein without notice. Before design-in or order placement, customers are advised to verify that the information is current on which orders or design activities are based. Cyrix warrants its products to conform to current specifications in accordance with Cyrix' standard warranty. Testing is performed to the extent necessary as determined by Cyrix to support this warranty. Unless explicitly specified by customer order requirements, and agreed to in writing by Cyrix, not all device characteristics are necessarily tested. Cyrix assumes no liability, unless specifically agreed to in writing, for customers' product design or infringement of patents or copyrights of third parties arising from use of Cyrix devices. No license, either express or implied, to Cyrix patents, copyrights, or other intellectual property rights pertaining to any machine or combination of Cyrix devices is hereby granted. Cyrix products are not intended for use in any medical, life saving, or life sustaining system. Information in this document is subject to change without notice.

May 22, 1997 9:39 am

C:\DATAOEM\!M2MXAP\103ap.fm5

Rev 1.6 New name, minor corrections

Rev 1.5 Corrected Page 11 0000h -> 00h

Rev 1.4 Corrected formatting and typos

Rev 1.3 Corrected Page 11, Changed page 9

Rev 1.2 Modified format

Rev 1.1 Table 1-1 was added, summarizing 6x86MX and 6x86 differences

Rev 1.0 First Release