

Floating-Point Arithmetic with the TMS32020

APPLICATION REPORT: SPRA011

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Floating-Point Arithmetic with the TMS32020

Abstract

This report presents algorithm and code implementing floating-point addition, subtraction, multiplication, and division with the TMS320. The support of floating-point operations by the TI processors has made possible some applications, such as implementation of the CCITT Adaptive Differential Pulse Code Modulation (ADPCM) algorithm and image/graphics operations.



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INTRODUCTION

The TMS32020 Digital Signal Processor is a fixed-point 16/32-bit microprocessor. However, it can also perform floating-point computations at a speed comparable to some dedicated floating-point processors.

The purpose of this application report is to analyze an implementation of floating-point addition, multiplication, and division on the TMS32020. The floating-point single-precision standard proposed by the IEEE will be examined. Using this standard, the TMS32020 performs a floating-point multiplication in 7.8 microseconds, a floating-point addition in 15.4 microseconds, and a floating-point division in 22.8 microseconds.

To illustrate floating-point formats and the tradeoffs involved in making a choice between different floating-point formats, a review of floating-point arithmetic notation and of addition, multiplication, and division algorithms is first presented.

FLOATING-POINT NOTATION

The floating-point number f may be written in floating-point format as

$$f = m \times b^e$$

where

m = mantissa
 b = base
 e = exponent

For example, 6,789,320 may be written as

$$0.6789320 \times 10^7$$

In this case,

$m = 0.6789320$
 $b = 10$
 $e = 7$

The two floating-point numbers f_1 and f_2 may be written as

$$f_1 = m_1 \times b^{e_1}$$

$$f_2 = m_2 \times b^{e_2}$$

Floating-point addition/subtraction, multiplication, and division for f_1 and f_2 are defined as follows:

$$f_1 \pm f_2 = (m_1 \pm m_2 \times b^{-(e_1 - e_2)}) \times b^{e_1} \text{ if } e_1 \geq e_2 \quad (1)$$

or

$$= (m_1 \times b^{-(e_2 - e_1)} \pm m_2) \times b^{e_2} \text{ if } e_1 < e_2$$

$$f_1 \times f_2 = m_1 \times m_2 \times b^{(e_1 + e_2)} \quad (2)$$

$$f_1/f_2 = (m_1/m_2) \times b^{(e_1 - e_2)} \quad (3)$$

A cursory examination of these expressions reveals some of the factors involved in the implementation of floating-point arithmetic. For addition, it is necessary to shift the mantissa of the floating-point number which has the smaller exponent to the right by the difference in the magnitude of the two exponents. This is shown in the multiplication by the terms

$$b^{-(e_1 - e_2)} \text{ and } b^{-(e_2 - e_1)}$$

This right shift can result in mantissa underflow. There are also possibilities for mantissa overflow. Addition and subtraction of exponents can lead to exponent overflow and overflow. To alleviate underflow and overflow, it is necessary to decide on some scheme for roundoff. For a detailed description and analysis of underflow and overflow conditions and rounding schemes, see reference 1.

It is desirable to have all numbers normalized, i.e., the mantissas of f_1 and f_2 have the most significant digit in the leftmost position. This provides the representation with the greatest accuracy possible for a fixed mantissa length. The result of any floating-point operation must also be normalized. The factors associated with normalization, overflow, and other characteristics of floating-point implementations are best illustrated with a few examples.

Consider the addition of two binary floating-point numbers f_1 and f_2 where

$$f_1 = 0.10100 \times 2^{011}$$

$$f_2 = 0.11100 \times 2^{001}$$

Both of these numbers are normalized, i.e., the first bit after the binary point is a 1. Addition requires equal exponents, so the fractions are aligned by shifting right the one with the smaller exponent and adjusting the smaller exponent. This yields

$$f_2 = 0.00111 \times 2^{011}$$

Then,

$$f_1 + f_2 = 0.10100 \times 2^{011} + 0.00111 \times 2^{011}$$

$$= 0.11011 \times 2^{011} = f_3$$

The sum may overflow the left end by one digit, thus requiring a postaddition adjustment or renormalization step. Since it is assumed that the register is only of a finite length, this renormalization will result in the loss of the lowest order bit.

Another example illustrates the overflow past the most significant bit. With an assumed register length of five, let

$$f_1 = 0.11100 \times 2^{011}$$

$$f_2 = 0.10101 \times 2^{001}$$

Then,

$$\begin{array}{r} 0.11100 \times 2^{011} = f_1 \\ + 0.0010101 \times 2^{011} = f_2 \\ \hline 1.0000101 \times 2^{011} = f_3 \end{array}$$

The significance of the two digits underlined in the right part of the mantissa is suspect, since it is assumed that the corresponding bits of f_1 are zero. The left underlined digit is the overflow past the most significant bit. To finish the addition, f_3 is shifted to the right and the exponent adjusted accordingly. Thus,

$$1.0000101 \times 2^{011} = f_3$$

The shift of the fraction and the adjustment of the exponent yield

$$0.10000101 \times 2^{100} = f_3$$

The result may be rounded, giving

$$0.10001 \times 2^{100} = f_3$$

or truncated, giving

$$0.10000 \times 2^{100} = f_3$$

FLOATING-POINT ALGORITHMS

Multiplication Algorithm

The algorithm for normalized floating-point multiplication is illustrated in Figure 1. This algorithm is an implementation of Equation 2 in the section on floating-point notation. The floating-point numbers being multiplied are A and B written as

$$A = m_A \times b^{e_A} \text{ and } B = m_B \times b^{e_B}$$

The result is

$$C = m_C \times b^{e_C}$$

For the resulting m_C , there are three special cases. The m_C may be zero, in which case there is a branch to Step 10 to set $C=0$. If $m_C \neq 0$, then the most significant bit will

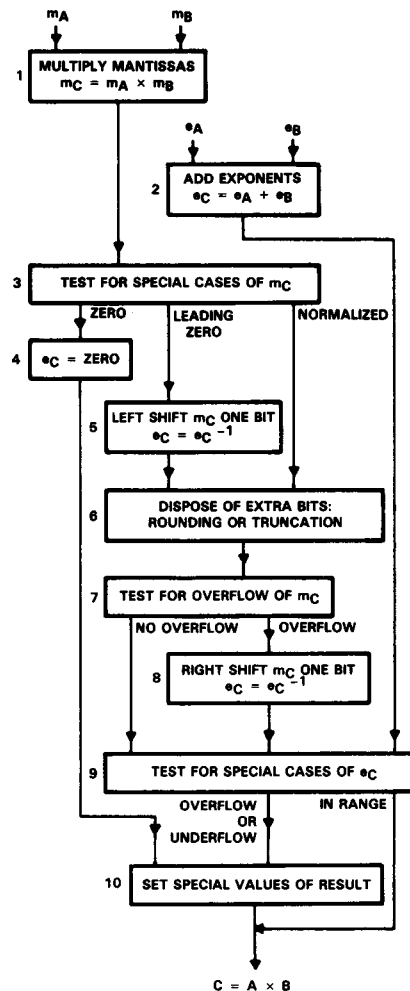


Figure 1. Floating-Point Multiplication

be in either the first or second leftmost bit. If the most significant bit is in the second leftmost bit, then a left shift of m_C is necessary (see Step 5). Otherwise, C is already in normalized form, and there is a branch to Step 6.

In Step 6, the desired rounding scheme is implemented. After this rounding, it is possible that m_C will overflow (see Step 7). In this case, it is necessary to right-shift m_C one bit (see Step 8). Special cases of e_C , are tested for in Step 9.

If there is an overflow or underflow of e_C , it is corrected in Step 10. Otherwise, the result is in range, and the calculation is complete.

Addition Algorithm

The implementation of normalized floating-point addition is more involved than for multiplication. This addition algorithm, outlined in Figure 2, is an implementation of Equation 1 in the section on floating-point notation.

In Step 1, e_A and e_B are compared to determine e_C . For this illustration of the algorithm, it is assumed that $e_A \leq e_B$. The right shift (d) required to align m_A is determined in Step 2. The procedure in Step 3 implements the right shift of m_A . In Step 4, the extra bits of m_A are discarded by using the desired rounding technique. The mantissas of A and B are then added in Step 5.

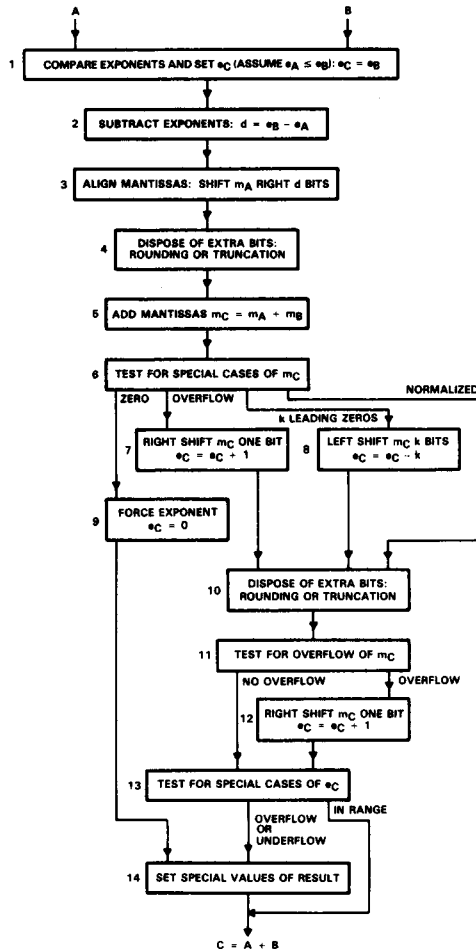


Figure 2. Floating-Point Addition

Now, the procedure becomes somewhat more involved. The m_C may be zero, in which case there is a branch to Step 9 which sets $e_C = 0$; a branch to Step 14 sets the special value of the result. The m_C may overflow, making a right shift of one necessary (see Step 7). The m_C may have k leading zeroes; therefore, a left shift of k is required. This normalization step is generally the most involved and time-consuming step to perform. The procedures in Steps 10, 11, and 12 round the m_C , test for a possible overflow due to the rounding, and adjust e_C accordingly. The special case of e_C is determined in Step 13. Finally, after Step 14, the sum $C = A + B$ is formed.

Division Algorithm

Floating-point division is more sophisticated than multiplication and addition since fixed-point processors such as the TMS32020 are not inherently capable of performing division. For example, $1/3 = 0.3333\dots$; only an approximation can be calculated since $1/3$ must be represented in a finite number of terms. Several algorithms can be implemented to find good approximations of such numbers. The algorithm implemented in this report is shown in Figure 3.

Step 1 shows the equivalent of A/B . In Step 2, the latter term is expanded using a power series of $1/(1 + X)$, where ϵ (BLO/BHI) is X (ϵ simply denotes that the term is right-shifted 16 bits forming the least significant bits of a 32-bit number). The third term in the power series only affects the LSB of a 32-bit result; therefore, this term and all the following terms can be dropped, as shown in Step 3.

The equation in Step 3 can be implemented on the TMS32020 in two steps. Assuming that the result is a 32-bit number Q and that it is composed of a 16-bit QHI and a 16-bit QLO , think of the equation in Step 3 in the following manner: $A/B = Q - \epsilon X$. The first term is a fair approximation of the result Q , and the second term is a correction term to obtain a better approximation. With this in mind, it can be shown that $(AHI + \epsilon ALO)/BHI$ will give a 16-bit quotient and a 16-bit remainder. Due to the architecture of the TMS32020, the 16-bit quotient will be in the low word of the accumulator and the remainder will be in the high word of the accumulator after the division. Since it is desirable

A divided by B

$$\text{where } A = AHI + \epsilon ALO$$

$$B = BHI + \epsilon BLO$$

$$\epsilon = \frac{1}{2^{\text{WORDSIZE}}} \cdot \frac{1}{2^{16}}$$

$$\text{STEP 1: } \frac{AHI + \epsilon ALO}{BHI + \epsilon BLO} = \frac{AHI + \epsilon ALO}{BHI} \left(\frac{1}{1 + \epsilon \left(\frac{BLO}{BHI} \right)} \right)$$

$$\text{STEP 2: } = \frac{AHI + \epsilon ALO}{BHI} \left(1 - \epsilon \left(\frac{BLO}{BHI} \right) + \epsilon^2 \left(\frac{BLO}{BHI} \right)^2 \dots \right)$$

$$\text{STEP 3: } = \frac{AHI + \epsilon ALO}{BHI} - \epsilon \left(\frac{BLO}{BHI} \right) \left(\frac{AHI + \epsilon ALO}{BHI} \right)$$

Figure 3. Division Equation

to have a floating-point result, the remainder must be divided by BHI to obtain the low word of the quotient. Now QHI and QLO have been calculated. When placing Q into the correction term (equation in Step 3), note that Q is equal to QHI + QLO. It can be shown that QLO will have no effect on the result since the correction term is multiplied by ϵ . Therefore, to calculate A divided by B, simply implement the following equation:

$$\frac{A}{B} = \frac{A}{BHI} - \epsilon \left(\frac{BLO}{BHI} \times QHI \right)$$

where the division is fixed binary (left-shifts and subtracts).

Figure 4 shows the implementation of the division algorithm that was outlined in Figure 3.

In Step 1, the dividend is right-shifted four times to prevent an overflow. Note that the result is not shifted left to compensate for this shift, because the normalization routine automatically does this. The shift causes the dividend to be limited to 27 significant bits instead of 31. In Step 2, a binary divide (left-shifts and subtracts) is implemented on the dividend by the high 16 bits of the divisor. The 32-bit result contains a quotient in the low 16 bits of the accumulator, and a remainder (R1) in the high 16 bits of the accumulator. R1 is left-shifted fifteen places in Step 3. The new R1 is divided by BHI in Step 4 to calculate the lower 16 bits of the quotient.

The quotient has now been approximated. The 32-bit result is composed of QHI and QLO, as shown in Figure 3. To obtain a better approximation, one term in the power series expansion must be added to the quotient. Therefore, the procedure in Step 5 calculates a 16-bit correction term, which is then added (or subtracted since it is the term following the "1" in the power series) to the 32-bit quotient.

Testing for an overflow of the resulting mantissa is necessary. Since the dividend was left-shifted four places, the resulting quotient will not be negative if an overflow occurred. To detect an overflow, bit 28 in the quotient must be tested. If this bit is a 1, an overflow occurred; if it is a 0, no overflow occurred. If an overflow has occurred, the exponent must be incremented. Finally, it is necessary to normalize the quotient and output the results.

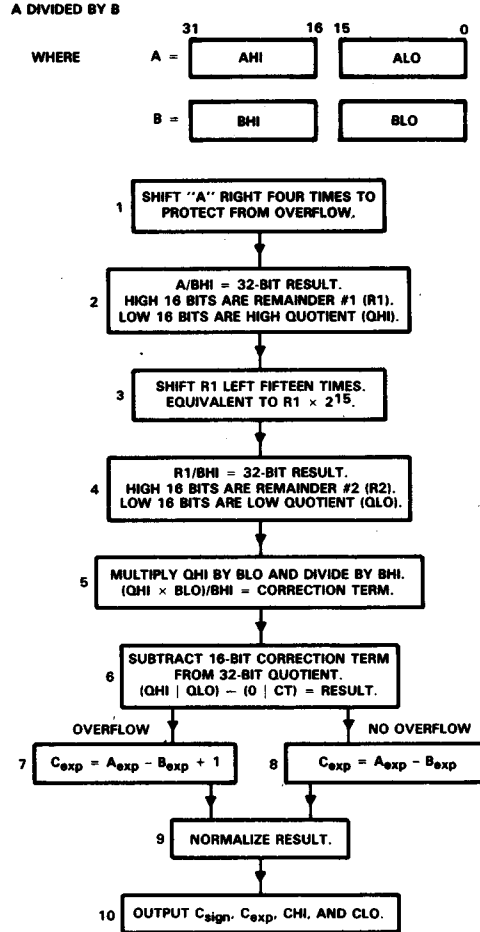


Figure 4. Floating-Point Division

IEEE FLOATING-POINT SINGLE-PRECISION FORMAT

Of interest is a set of formats known as the IEEE standard. This IEEE recommended format consists of a variety of precision formats (single, double, single-extended, and double-extended). The IEEE has also proposed several techniques for handling special cases such as overflow, underflow, $\pm \infty$, and rounding. For complete details, the reader is referred to the proposed IEEE standard.²

The single-precision format is a 32-bit format consisting of a 1-bit sign field *s*, an 8-bit biased exponent *e*, and a 23-bit fraction *f* (see Figure 5). The value of a binary floating-point number *X* is determined as follows:

$$X = (-1)^s \times 2^{(e-127)} \times 1.f$$

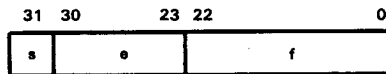


Figure 5. IEEE Floating-Point Single-Precision Format

The advantage of this format is that it is structured in such a way as to provide easy storage and straightforward input/output operations on 8-, 16- and 32-bit processors. The disadvantage with this format is that the large mantissa will generally span several words of memory.

FLOATING-POINT IMPLEMENTATION

IEEE Implementation

The IEEE single-precision format is described here as it applies to the addition, multiplication, and division algorithms. In these floating-point routines written for the TMS32020, all results are truncated to 31 bits to provide more flexibility in the user's development of a rounding scheme suitable for his application. The representations of $\pm \infty$ are ignored so that the user can decide how to handle these exceptions in a manner that is appropriate for his particular application.

I/O Considerations

The first consideration is the internal representation of the binary floating-point number. If the number is read into the TMS32020 as two 16-bit words, some processing is then necessary to put the floating-point number into a representation which is easier to process. The representation used in the TMS32020 programs in the appendices is shown in Figure 6. This internal representation may be arrived at by a simple manipulation of the IEEE bit fields. For this particular algorithm, it is assumed that the floating-point number is input to the TMS32020 as the four 16-bit fields shown in Figure 6. However, the user can easily supply his

own routine to arrive at this format from two 16-bit inputs to the TMS32020 where the inputs contain the IEEE single-precision format.

The format in Figure 6 was chosen to minimize the execution time of the floating-point addition, multiplication, and division routines. The format of the result is shown in Figure 7. Notice that it is identical to the format in Figure 5 except for CLO. CLO has its 16 most significant bits valid for both the addition, multiplication, and division routines.

Normalization

Since the floating-point routines require normalization, a partial binary search algorithm is implemented in the addition and division routines in the appendices. To begin the normalization routine, note that all mantissas can be considered to be positive with the format used for the result shown in Figure 7. The binary search for the most significant bit (the leftmost 1 since the mantissa is positive) is illustrated in Figure 8.

The first move is to split the result into CHI and CLO. If CHI $\neq 0$, the most significant bit (MSB) is the CHI; otherwise, it is the CLO. For this example, it is in CLO.

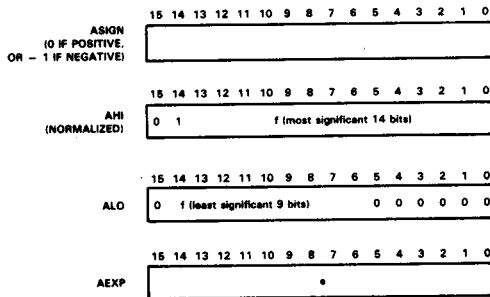


Figure 6. Floating-Point Representation

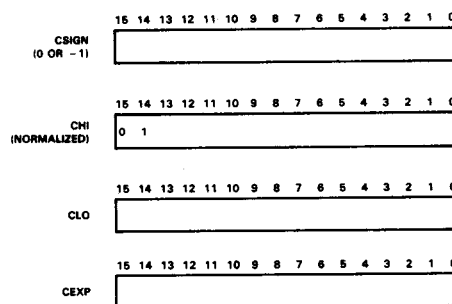


Figure 7. Result Representation

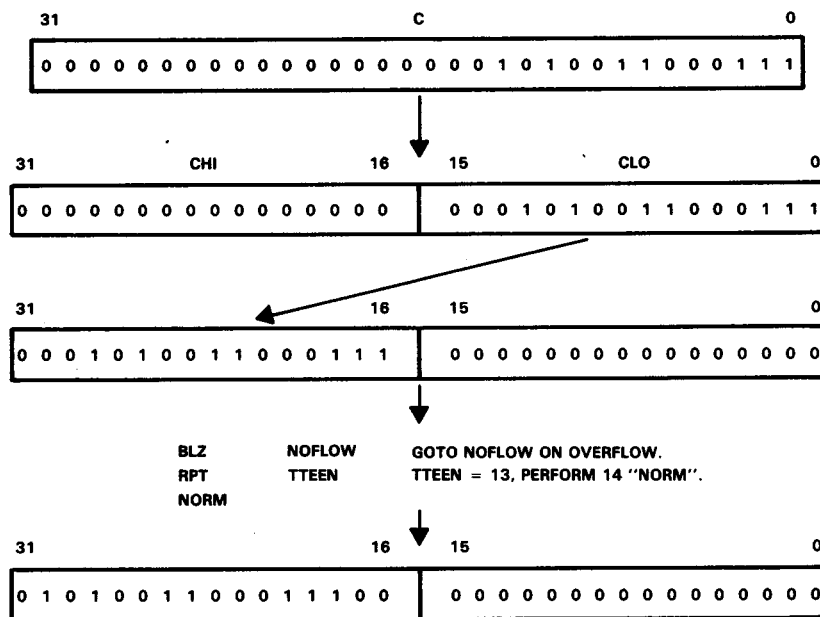


Figure 8. Partial Binary Search

The next step is to form a 32-bit result with CLO in the most significant word position. It is now possible for the MSB to be in the highest bit location since CLO has been left-shifted 16 times. If this is the case, an overflow has occurred, and the result must be right-shifted once. The normalization routine tests this by branching to NOFLOW if the result is negative. If the number is not negative, the normalization can continue.

The NORM instruction is used in the repeat mode to complete the normalization. Note that this whole normalization routine can be replaced by the following two instructions: RPTK 29 and NORM. The RPTK instruction causes the NORM instruction to be repeated 30 times, thus normalizing a 32-bit number. This method is not implemented here due to the timing. These two instructions always take 31 cycles to normalize a 32-bit number. The normalization routine here takes only 22 cycles (worst case) for normalizing a 32-bit number. Therefore, if program space is more important than timing efficiency, it is best to replace the normalization routine with these two instructions.

Added Precision

As illustrated in Figure 7, the 16 most significant bits of CLO are valid, i.e., C is valid for 31 places beyond the

binary point. Oftentimes the user is not as concerned with the IEEE standard as in being certain that he has enough accuracy for his particular application. Since the TMS32020 uses 16-bit words, the routines in the appendices implicitly maintain a 30-bit mantissa. They also implicitly use a 16-bit exponent. If the user desires this added accuracy and dynamic range, then it is readily implementable with no additional cost in execution time. The normalization for the addition, as mentioned previously, operates over the entire 32-bit accumulator. For the strict IEEE format, the user will only want to normalize over the 25 most significant bits of the accumulator. The structure of the normalization routine makes this modification simple.

The routines in the appendices make no provision for the representation of $\pm \infty$ and exponent underflow and overflow. The user of the routines should consider the degree of significance of these results and the way they should be handled for his particular application. Since these routines are written to operate at maximum speed, truncation of results is used. If the user desires to implement a rounding scheme, then he will also need to check for the possibility of overflow due to the rounding scheme. This step is shown in the multiplication, addition, and division flowcharts (see Figures 1, 2, and 3).

SUMMARY

The TMS32020 may be used to perform floating-point operations with great accuracy, wide dynamic range, and high-speed execution. The design engineer has the responsibility of deciding what type of floating-point format is best for his application. To aid in understanding floating-point operations, several examples have been given that illustrate the manipulations necessary to implement floating-point addition, multiplication, and division algorithms. Flowcharts for these algorithms are also included. The appendices contain the TMS32020 code for the IEEE floating-point single-precision format used in addition, multiplication, and division. The addition and multiplication routines may also be used without modification to implement a format with up to a 30-bit mantissa and a 16-bit exponent without any increase in execution time.

ACKNOWLEDGEMENTS

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1. D.J. Kuck, *The Structure of Computers and Computations*, Vol 1, John Wiley & Sons (1978).
2. J. Coonen et al, "A Proposed Standard for Binary Floating-Point Arithmetic," *ACM Signum Newsletter*, 4-12 (October 1979).
3. Donald E. Knuth, *Seminumerical Algorithms*, Vol 2, 2nd Edition, Addison-Wesley (1981).

APPENDIX A

FLTADD 32020 FAMILY MACRO ASSEMBLER PC 1.0 85.157 11:47:00 08-19-86
PAGE 0001

```

0001      *
0002      *
0003      *   THIS IS A FLOATING-POINT ADDITION ROUTINE WHICH   *
0004      *   IMPLEMENTS THE IEEE PROPOSED FLOATING-POINT   *
0005      *   FORMAT ON THE TMS32020.                         *
0006      *
0007      *
0008      *
0009      *   INITIAL FORMAT (ALL 16 BIT WORDS)
0010      *   -----
0011      *   ;   ALL 0 OR 1   ;           ASIGN (0 OR -1)
0012      *   -----
0013      *
0014      *
0015      *   -----
0016      *   ;0|.  15 BITS  ;           AHI (NORMALIZED)
0017      *   -----
0018      *
0019      *   -----
0020      *   ;0|  9 BITS |--0-|         ALO
0021      *   -----
0022      *
0023      *   -----
0024      *   ;           ;           AEXP (-127 TO 128)
0025      *   -----
0026      *
0027      *   TO CORRESPOND WITH IEEE FORMAT,
0028      *   INPUT 0.1F * 2 ** (E + 1)
0029      *   INSTEAD OF 1.F * 2 **E, AND SUBTRACT 127 FROM E.
0030      *
0031      *   THE FINAL FORMAT IS THE SAME AS THE INITIAL FORMAT
0032      *   EXCEPT THAT FOR CLO WE HAVE:
0033      *
0034      *   -----
0035      *   ;   16 BITS   ;           CLO
0036      *   -----
0037      *
0038      *   ALL 16 BITS OF CLO ARE VALID. ANYTHING PAST THESE HAS
0039      *   BEEN TRUNCATED.
0040      *
0041      *
0042      *
0043      *   WORST CASE (EXCLUDING INITIALIZATION AND I/O):
0044      *   15.4 MICROSECONDS.
0045      *   THIS TIMING INCLUDES THE NORMALIZATION.
0046      *   WORDS OF PROGRAM MEMORY: 217
0047      *
0048      *
0049      *
0050      *   IDT      ' FLTADD'
0051      *   0000  AORG
0052      *   0001  ASIGN  EQU   0
0053      *   0002  AEXP   EQU   1
0054      *   0003  AHI    EQU   2
0055      *   0004  ALO    EQU   3
0056      *   0005  BSIGN EQU   4
0057      *   0006  BEXP   EQU   5
0058      *   0007  BH!   EQU   6

```

```

0058 0007 BLO EQU 7
0059 0008 CSIGN EQU 8
0060 0009 CEXP EQU 9
0061 000A CHI EQU 10
0062 000B CLO EQU 11
0063 000C D EQU 12
0064 000D ONE EQU 13
0065 000E TEMP EQU 14
0066 000F THREE EQU 15
0067 0010 SIXT EQU 16
0068 0011 RESID EQU 17
0069 0012 TTEEN EQU 18
0070 *
0071 * INITIALIZATION
0072 *
0073 0000 C804 LDPK 4 BEGIN ON PAGE 4.
0074 0001 CE07 SSXM SET SIGN EXTENSION.
0075 0002 5589 LARP 1
0076 0003 D100 LRLK AR1,>200
0004 0200
0077 0005 CB07 RPTK 7
0078 0006 80A0 IN *+,PA0
0079 0007 5588 LARP 0
0080 0008 C000 LARK AR0,0 CLEAR EXPONENT REGISTER.
0081 0009 CA01 LACK 1
0082 000A 6000 SACL ONE ONE = 1
0083 000B CA10 LACK 16
0084 000C 6010 SACL SIXT
0085 000D CA03 LACK 3
0086 000E 600F SACL THREE
0087 000F CA0D LACK 13
0088 0010 6012 SACL TTEEN
0089 *
0090 * BEGIN FLOATING POINT ADD
0091 *
0092 0011 2001 UP LAC AEXP FIND LARGEST NUMBER.
0093 0012 1005 SUB BEXP
0094 0013 F680 BZ AEQB IF EXP ARE THE SAME, JUMP TO AEQB.
0014 0043
0095 0015 F380 BLZ ALTB IF A IS LESS THAN B, JUMP TO ALTB.
0016 004D
0096 *
0097 0017 CE23 AGTB NEG
0098 0018 0010 ADD SIXT D = (16-D)
0099 0019 F380 BLZ AI JUMP IF EXP DIFFERENCE IS > 16
001A 0028
0100 *
0101 * EXPONENT DIFFERENCE < 16
0102 *
0103 001B 600C SACL D
0104 001C 3C0C LT D
0105 001D 4206 LACT BHI BHI IS SHIFTED RIGHT "D" TIMES.
0106 001E 6806 SACH BHI
0107 001F 6011 SACL RESID RESIDUAL BITS MUST BE MAINTAINED.
0108 0020 4207 LACT BLO BLO IS SHIFTED RIGHT "D" TIMES.
0109 0021 CE18 SFL MSB (THE 0) IS SHIFTED AWAY.
0110 0022 6807 SACH BLO

```

```

0111 0023 2007      LAC   BLO
0112 0024 4D11      OR    RESID   GET BITS THAT WERE SHIFTED FROM BHI.
0113 0025 6007      SACL  BLO
0114 0026 FF80      B     A2
      0027 0031
0115                *
0116                *   EXPONENT DIFFERENCE >16
0117                *
0118 0028 0010  A1   ADD   SIXT
0119 0029 F380      BLZ   A3           JUMP IF EXPONENT DIFF > 32
      002A 0039
0120 002B 600C      SACL  D
0121 002C 3C0C      LT    D
0122 002D 4206      LACT  BHI
0123 002E 6807      SACH  BLO
0124 002F CA00      ZAC
0125 0030 6006      SACL  BHI
0126 0031 2000  A2   LAC   ASIGN   A IS LARGER THAN B.
0127 0032 6008      SACL  CSIGN   THEREFORE, CSIGN = ASIGN.
0128 0033 2001      LAC   AEXP   ALIGN THE B MANTISSA.
0129 0034 6009      SACL  CEXP
0130 0035 2103      LAC   ALO,1   GET RID OF EXTRA BIT.
0131 0036 6003      SACL  ALO
0132 0037 FF80      B     CHKSGN  DO BOTH NUMBERS HAVE THE SAME SIGN?
      0038 0078
0133                *
0134                *   A >> B , RESULT = A
0135                *
0136 0039 2002  A3   LAC   AHI
0137 003A 600A      SACL  CHI
0138 003B 2103      LAC   ALO,1
0139 003C 600B      SACL  CLO           A IS LARGER THAN B
0140 003D 2000      LAC   ASIGN   THEREFORE CSIGN = ASIGN
0141 003E 6008      SACL  CSIGN
0142 003F 2001      LAC   AEXP
0143 0040 6009      SACL  CEXP
0144 0041 FF80      B     AROUND
      0042 00D6
0145                *
0146 0043 2000  AEQB LAC   ASIGN   IF SIGNS ARE THE SAME, CSIGN = ASIGN
0147 0044 6008      SACL  CSIGN
0148 0045 2103      LAC   ALO,1   ALIGN MANTISSAS.
0149 0046 6003      SACL  ALO
0150 0047 2107      LAC   BLO,1
0151 0048 6007      SACL  BLO
0152 0049 2001      LAC   AEXP   SET C EXPONENT = A EXPONENT.
0153 004A 6009      SACL  CEXP
0154 004B FF80      B     CHKSGN  DO BOTH NUMBERS HAVE THE SAME SIGN?
      004C 0078
0155                *
0156 004D 0010  ALTB ADD   SIXT   D = (16-D)
0157 004E F380      BLZ   B1           JUMP IF EXP DIFF > 16
      004F 005D
0158 0050 600C      SACL  D
0159 0051 3C0C      LT    D
0160 0052 4202      LACT  AHI   AHI GETS SHIFTED "D" TIMES.
0161 0053 6802      SACL  AHI

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0162 0054 6011      SACL  RESID  MAINTAIN EXTRA BITS.
0163 0055 4203      LACT  ALO    ALO GETS SHIFTED "D" TIMES.
0164 0056 CE18      SFL    MSB (THE 0) IS SHIFTED AWAY.
0165 0057 6803      SACH  ALO
0166 0058 2003      LAC   ALO
0167 0059 4011      OR    RESID  GET RESIDUAL BITS.
0168 005A 6003      SACL  ALO
0169 005B FF80      B     B2
      005C 0066
0170      *
0171      *   EXPONENT DIFFERENCE > 16
0172      *
0173 005D 0010  B1  ADD   SIXT
0174 005E F380      BLZ   B3      JUMP IF ,EXP DIFF > 32
      005F 006E
0175 0060 600C      SACL  D
0176 0061 3C0C      LT   D
0177 0062 4202      LACT  AHI
0178 0063 6803      SACH  ALO
0179 0064 CA00      ZAC
0180 0065 6002      SACL  AHI
0181 0066 2004  B2  LAC   BSIGN  B IS THE BIGGEST NUMBER.
0182 0067 6008      SACL  CSIGN  THEREFORE, LET THE SIGN OF C=BSIGN.
0183 0068 2005      LAC   BEXP   SET C EXPONENT = B EXPONENT.
0184 0069 6009      SACL  CEXP
0185 006A 2107      LAC   BLO,1  GET RID OF EXTRA BIT.
0186 006B 6007      SACL  BLO
0187 006C FF80      B     CHKSGN  DO BOTH NUMBERS HAVE THE SAME SIGN?
      006D 0078
0188      *
0189      *   B >> A ,   RESULT = B
0190      *
0191 006E 2006  B3  LAC   BHI
0192 006F 600A      SACL  CHI
0193 0070 2107      LAC   BLO,1
0194 0071 600B      SACL  CLO    B IS THE BIGGEST NUMBER
0195 0072 2004      LAC   BSIGN  THEREFORE, LET THE SIGN OF C=BSIGN
0196 0073 6008      SACL  CSIGN
0197 0074 2005      LAC   BEXP   SET C EXPONENT = B EXPONENT
0198 0075 6009      SACL  CEXP
0199 0076 FF80      B     AROUND
      0077 00D6
0200      *
0201 0078 2000  CHKSGN LAC   ASIGN  CHECK THE SIGNS.
0202 0079 1004      SUB   BSIGN
0203 007A F680      BZ   ADNOW   IF THEY ARE THE SAME, JUST ADD.
      007B 00A9
0204 007C F380      BLZ   AISNEG
      007D 00BC
0205 007E 4002  BISNEG ZALH  AHI    DO (|A| - |B|),
0206 007F 4903      ADDS  ALO    SINCE B < 0 AND A > 0.
0207 0080 4507      SUBS  BLO
0208 0081 4406      SUBH  BHI
0209 0082 F680      BZ   CZERO
      0083 009A
0210 0084 F380      BLZ   CNEG
      0085 00A1
    
```

0211	0086	680A	SACH	CHI	
0212	0087	600B	SACL	CLO	
0213	0088	CA00	ZAC		
0214	0089	6008	SACL	CSIGN	
0215	008A	FF80	B	NORMAL	GO AND NORMALIZE RESULT.
	008B	00B3			
0216	008C	4006	AISNEG ZALH	BHI	DO (!B! - !A!),
0217	008D	4907	ADDS	BLO	SINCE A < 0 AND B > 0.
0218	008E	4503	SUBS	ALO	
0219	008F	4402	SUBH	AHI	
0220	0090	F680	BZ	CZERO	
	0091	009A			
0221	0092	F380	BLZ	CNEG	
	0093	00A1			
0222	0094	680A	SACH	CHI	
0223	0095	600B	SACL	CLO	
0224	0096	CA00	ZAC		
0225	0097	6008	SACL	CSIGN	
0226	0098	FF80	B	NORMAL	GO AND NORMALIZE RESULTS.
	0099	00B3			
0227			*		
0228	009A	CA00	CZERO ZAC		HERE, ONLY IF RESULT = 0.
0229	009B	6009	SACL	CEXP	
0230	009C	6008	SACL	CSIGN	
0231	009D	600A	SACL	CHI	
0232	009E	600B	SACL	CLO	
0233	009F	FF80	B	AROUND	OUTPUT A ZERO.
	00A0	00D6			
0234			*		
0235	00A1	CE1B	CNEG ABS		HERE, IF RESULT IS NEGATIVE.
0236	00A2	680A	SACH	CHI	
0237	00A3	600B	SACL	CLO	
0238	00A4	D001	LALK	>FFFF	
	00A5	FFFF			
0239	00A6	6008	SACL	CSIGN	
0240	00A7	FF80	B	NORMAL	GO NORMALIZE RESULT.
	00A8	00B3			
0241			*		
0242	00A9	4002	ADNOW ZALH	AHI	IF SIGNS ARE THE SAME, JUST ADD.
0243	00AA	4903	ADDS	ALO	
0244	00AB	4907	ADDS	BLO	
0245	00AC	4806	ADDH	BHI	
0246	00AD	680A	SACH	CHI	
0247	00AE	600B	SACL	CLO	
0248	00AF	F080	BV	OVFLOW	DID AN OVERFLOW OCCUR?
	00B0	00C4			
0249	00B1	F680	BZ	CZERO	IS RESULT = 0 ?
	00B2	009A			
0250			*		
0251			*	NORMALIZE	
0252			*		
0253	00B3	200A	NORMAL LAC	CHI	DOES CHI HAVE THE MSB?
0254	00B4	F680	BZ	LOI	
	00B5	00BC			
0255	00B6	400A	ZALH	CHI	IF YES, NORMALIZE RESULT.
0256	00B7	490B	ADDS	CLO	
0257	00B8	4B12	RPT	TTEEN	WILL PERFORM 14 "NORMS"

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0258 00B9 CEA2      NORM
0259 00BA FF80      B      OUTPUT      GO OUTPUT RESULTS.
      00BB 00D0
0260 00BC 400B     LOI      ZALH      CLO      HERE IF CLO HAS MSB.
0261 00BD C010      LARK      ARO,16  OFFSET EXPONENT BY 16.
0262 00BE F380      BLZ      NOFLOW     DID BIT SEARCH CAUSE OVERFLOW?
      00BF 00CD
0263 00C0 4B12      RPT      TTEEN     IF NOT, NORMALIZE RESULT.
0264 00C1 CEA2      NORM
0265 00C2 FF80      B      OUTPUT      GO OUTPUT RESULT.
      00C3 00D0
0266 *
0267 *
0268 *      FINISHED WITH NORMALIZATION
0269 *
0270 *      HERE ONLY IF OVERFLOW OCCURRED DURING ADDITION
0271 *
0272 *
0273 00C4 CE06     OVFLOW  RSXM      RESET SIGN EXTENSION TO SHIFT RIGHT
0274 00C5 CE19      SFR      SHIFT RIGHT.
0275 00C6 680A     SACH     CHI      STORE NORMALIZED MANTISSA.
0276 00C7 600B     SACL     CLO
0277 00C8 2009     LAC     CEXP     DECREMENT EXPONENT.
0278 00C9 000D     ADD     ONE
0279 00CA 6009     SACL     CEXP
0280 00CB FF80      B      AROUND     GO OUTPUT RESULTS.
      00CC 00D6
0281 *
0282 *      OVERFLOW OCCURRED DURING BIT SEARCH
0283 *
0284 00CD 5590     NOFLOW  MAR      *-      DECREMENT EXPONENT.
0285 00CE CE06     RSXM
0286 00CF CE19     SFR      RSXM FOR LOGICAL RIGHT SHIFT.
      *      PERFORM RIGHT SHIFT.
0287 *
0288 *
0289 *      TAKE CARE OF EXPONENT & NORMALIZED MANTISSA,
0290 *      THEN OUTPUT RESULTS.
0291 *
0292 *
0293 00D0 700E     OUTPUT  SAR      ARO,TEMP  HERE AFTER NORMALIZATION.
0294 00D1 680A     SACH     CHI      SAVE NORMALIZED MANTISSA.
0295 00D2 600B     SACL     CLO
0296 00D3 2009     LAC     CEXP     ADJUST EXPONENT.
0297 00D4 100E     SUB     .TEMP
0298 00D5 6009     SACL     CEXP
0299 *
0300 00D6 5589     AROUND  LARP     1      RESET POINTER.
0301 00D7 4B0F     RPT     THREE
0302 00D8 E0A0     OUT     *+,PA0
0303 00D9 CE1F     IDLE
NO ERRORS, NO WARNINGS

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APPENDIX B

NO\$IDT 32020 FAMILY MACRO ASSEMBLER PC0.7 84.348 15:24:53 03-27-85
 *** PRERELEASE ***

PAGE 0001

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0001            *****
0002            *
0003            *    THIS IS A FLOATING-POINT MULTIPLICATION ROUTINE WHICH *
0004            *    IMPLEMENTS THE IEEE PROPOSED FLOATING-POINT FORMAT *
0005            *    ON THE TMS32020.                                       *
0006            *                                                               *
0007            *****
0008            *
0009            *    INITIAL FORMAT (ALL 16-BIT WORDS)
0010            *    -----
0011            *    |    ALL 0 OR 1    |            ASIGN (0 OR -1)
0012            *    -----
0013            *
0014            *    -----
0015            *    |0|.    15 BITS    |            AHI (NORMALIZED)
0016            *    -----
0017            *
0018            *    -----
0019            *    |0|.    9 BITS |--0-|            ALO
0020            *    -----
0021            *
0022            *    -----
0023            *    |                    |            AEXP (-127 TO 128)
0024            *    -----
0025            *
0026            *    TO CORRESPOND WITH IEEE FORMAT,
0027            *    INPUT 0.1F * 2 ** (E + 1)
0028            *    INSTEAD OF 1.F * 2 **E, AND SUBTRACT 127 FROM E.
0029            *
0030            *    THE FINAL FORMAT IS THE SAME AS THE INITIAL FORMAT
0031            *    EXCEPT THAT FOR CLO WE HAVE:
0032            *
0033            *    -----
0034            *    |    16 BITS    |            CLO
0035            *    -----
0036            *
0037            *    ALL 16 BITS OF CLO ARE VALID. ANYTHING PAST THESE HAS
0038            *    BEEN TRUNCATED.
0039            *
0040            *****
0041            *
0042            *    WORST CASE (EXCLUDING INITIALIZATION AND I/O):    *
0043            *    7.8 MICROSECONDS.                                       *
0044            *    THIS TIMING INCLUDES THE NORMALIZATION.               *
0045            *    WORDS OF PROGRAM MEMORY: 60                               *
0046            *                                                               *
0047            *****
0048            *
0049            *    AORG
0050            *    0000    ASIGN    EQU    0
0051            *    0001    AEXP    EQU    1
0052            *    0002    AHI     EQU    2
0053            *    0003    ALO     EQU    3
0054            *    0004    BSIGN   EQU    4
0055            *    0005    BEXP    EQU    5
0056            *    0006    BHI     EQU    6
  
```

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0057 0007 BLO EQU 7
0058 0008 CSIGN EQU 8
0059 0009 CEXP EQU 9
0060 000A CHI EQU 10
0061 000B CLO EQU 11
0062 000C THI EQU 12
0063 000D NEGONE EQU 13
0064 000E TLO EQU 14
0065 000F TEMP EQU 15
0066 *
0067 *
0068 * INITIALIZATION
0069 *
0070 *
0071 *
0072 0000 C804 LDPK 4 BEGIN ON PAGE 4.
0073 0001 CE07 SSXM SET SIGN EXTENSION.
0074 0002 5589 LARP 1
0075 0003 D100 LRLK AR1,>200
0004 0200
0076 0005 CB07 RPTK 7 READ NUMBERS INTO BLOCK B0.
0077 0006 80A0 IN ** ,PA0
0078 0007 C000 LARK AR0,0 CLEAR EXPONENT REGISTER.
0079 0008 5588 LARP 0
0080 0009 D001 LALK >FFFF
000A FFFF
0081 000B 600D SACL NEGONE NEGONE = -1
0082 *
0083 *
0084 * BEGIN FLOATING-POINT MULTIPLICATION.
0085 *
0086 *
0087 000C 2001 UP LAC AEXP ADD EXPONENTS.
0088 000D 0005 ADD BEXP
0089 000E 6009 SACL CEXP
0090 *
0091 000F 3C03 LT ALO FIRST PRODUCT (ALO * BHI)
0092 0010 3806 MPY BHI
0093 0011 CE14 PAC
0094 0012 680C SACH THI
0095 0013 600E SACL TLO
0096 *
0097 0014 3C02 LT AHI SECOND PRODUCT (AHI * BLO)
0098 0015 3807 MPY BLO
0099 *
0100 0016 CE15 AFAC HAS EFFECT OF (AHI * BLO + ALO * BHI) * 2 ** -15.
0101 0017 CE15 AFAC
0102 *
0103 0018 480C ADDH THI
0104 0019 490E ADDS TLO
0105 001A 680C SACH THI
0106 *
0107 001B 3806 MPY BHI (AHI * BHI)
0108 001C CE14 PAC
0109 001D 490C ADDS THI
0110 *
  
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PAGE 0003

```
0111 001E 690A      SACH  CHI,1  GET RID OF EXTRA SIGN BITS.
0112 001F 610B      SACL  CLO,1
0113                *
0114 0020 F580      BNZ   OK    IS RESULT ZERO?
        0021 0026
0115 0022 CA00      ZAC
0116 0023 6009      SACL  CEXP
0117 0024 FF80      B     SETSIN
        0025 002F
0118                *
0119 0026 400A      OK    ZALH  CHI  NORMALIZE AND WRAP UP.
0120 0027 490B      ADDS  CLO
0121 0028 CEA2      NORM
0122 0029 680A      SACH  CHI
0123 002A 600B      SACL  CLO
0124 002B 700F      SAR   ARO,TEMP
0125 002C 2009      LAC   CEXP
0126 002D 100F      SUB   TEMP
0127 002E 6009      SACL  CEXP
0128                *
0129 002F 4100      SETSIN ZALS  ASIGN  WHAT IS SIGN OF RESULT?
0130 0030 4C04      XOR   BSIGN
0131 0031 F580      BNZ   NEG
        0032 0037
0132 0033 CA00      ZAC
0133 0034 6008      SACL  CSIGN
0134 0035 FF80      B     OUTPUT
        0036 0039
0135 0037 200D      NEG   LAC   NEGONE
0136 0038 6008      SACL  CSIGN
0137 0039 5589      OUTPUT LARP  1  OUTPUT RESULTS.
0138 003A CB03      RPTK  3
0139 003B E0A0      OUT   **,,PA0
0140 003C CE1F      IDLE
NO ERRORS, NO WARNINGS
```

APPENDIX C

NO#IDT 32020 FAMILY MACRO ASSEMBLER PC0.7 84.348 15:25:17 03-27-85
 *** PRERELEASE ***

PAGE 0001

```

0001                    *****
0002                    *
0003                    *    THIS IS A FLOATING-POINT DIVISION ROUTINE WHICH    *
0004                    *    IMPLEMENTS THE IEEE PROPOSED FLOATING-POINT FORMAT    *
0005                    *    ON THE TMS32020.                                    *
0006                    *                                                            *
0007                    *****
0008                    *
0009                    *    INITIAL FORMAT (ALL 16-BIT WORDS)                    *
0010                    *    -----                                                *
0011                    *    |    ALL 0 OR 1    |            ASIGN (0 OR -1)            *
0012                    *    -----                                                *
0013                    *                                                            *
0014                    *    -----                                                *
0015                    *    |0|.    15 BITS    |            AHI (NORMALIZED)            *
0016                    *    -----                                                *
0017                    *                                                            *
0018                    *    -----                                                *
0019                    *    |0|    9 BITS |--0-|            ALO                    *
0020                    *    -----                                                *
0021                    *                                                            *
0022                    *    -----                                                *
0023                    *    |                                    |            AEXP (-127 TO 128)            *
0024                    *    -----                                                *
0025                    *                                                            *
0026                    *    TO CORRESPOND WITH IEEE FORMAT,                    *
0027                    *    INPUT 0.1F * 2 ** (E + 1)                            *
0028                    *    INSTEAD OF 1.F * 2 **E, AND SUBTRACT 127 FROM E.            *
0029                    *                                                            *
0030                    *    THE FINAL FORMAT IS THE SAME AS THE INITIAL FORMAT            *
0031                    *    EXCEPT THAT FOR CLO WE HAVE:                            *
0032                    *                                                            *
0033                    *    -----                                                *
0034                    *    |    16 BITS    |            CLO                    *
0035                    *    -----                                                *
0036                    *                                                            *
0037                    *    ALL 16 BITS OF CLO ARE VALID. ANYTHING PAST THESE HAS            *
0038                    *    BEEN TRUNCATED.                                            *
0039                    *                                                            *
0040                    *****
0041                    *                                                            *
0042                    *    WORST CASE (EXCLUDING INITIALIZATION AND I/O):            *
0043                    *    22.8 MICROSECONDS.                                    *
0044                    *    THIS TIMING INCLUDES THE NORMALIZATION.                    *
0045                    *    WORDS OF PROGRAM MEMORY: 92                            *
0046                    *                                                            *
0047                    *****
0048                    *
0049                    *    AORG    0
0050                    *    0000    ASIGN    EQU    0
0051                    *    0001    AEXP    EQU    1
0052                    *    0002    AHI     EQU    2
0053                    *    0003    ALO     EQU    3
0054                    *    0004    BSIGN   EQU    4
0055                    *    0005    BEXP    EQU    5
0056                    *    0006    BHI     EQU    6
  
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0057	0007	BLO	EQU	7	
0058	0008	CSIGN	EQU	8	
0059	0009	CEXP	EQU	9	
0060	000A	CHI	EQU	10	
0061	000B	CLO	EQU	11	
0062	000C	NEGONE	EQU	12	
0063	000D	TEMP	EQU	13	
0064	000E	FOUR	EQU	14	
0065	000F	QM	EQU	15	
0066	0010	QL	EQU	16	
0067	0011	R1	EQU	17	
0068	0012	R2	EQU	18	
0069	0013	CL	EQU	19	
0070	0014	M1000	EQU	20	
0071	0015	ONE	EQU	21	
0072	0016	THREE	EQU	22	
0073	0017	FIFTEEN	EQU	23	
0074	0018	THIRTY	EQU	24	
0075	0019	TTEEN	EQU	25	
0076		*			
0077		*			
0078		*	INITIALIZATION		
0079		*			
0080		*			
0081		*			
0082	0000	C804	LDPK	4	BEGIN ON PAGE 4.
0083	0001	CE07	SSXM		SET SIGN EXTENSION.
0084	0002	5589	LARP	1	
0085	0003	D100	LRLK	AR1,>200	
	0004	0200			
0086	0005	CB07	RPTK	7	READ NUMBERS INTO BLOCK B0.
0087	0006	80A0	IN	** ,PA0	
0088	0007	5588	LARP	0	
0089	0008	C000	LARK	AR0,0	CLEAR EXPONENT REGISTER.
0090	0009	D001	LALK	>FFFF	
	000A	FFFF			
0091	000B	600C	SACL	NEGONE	NEGONE = -1
0092	000C	D001	LALK	>1000	
	000D	1000			
0093	000E	6014	SACL	M1000	M1000 = >1000
0094	000F	CA04	LACK	4	
0095	0010	600E	SACL	FOUR	FOUR = 4
0096	0011	CA01	LACK	1	
0097	0012	6015	SACL	ONE	ONE = 1
0098	0013	CA03	LACK	3	
0099	0014	6016	SACL	THREE	THREE = 3
0100	0015	CA0F	LACK	15	
0101	0016	6017	SACL	FIFTEEN	FIFTEEN = 15
0102	0017	CA1E	LACK	30	
0103	0018	6018	SACL	THIRTY	THIRTY = 30
0104	0019	CA0D	LACK	13	
0105	001A	6019	SACL	TTEEN	TTEEN = 13
0106	001B	CA00	ZAC		
0107	001C	6009	SACL	CEXP	CLEAR CEXP
0108		*			
0109		*			


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0110          *      FINISHED WITH INITIALIZATION
0111          *
0112          *
0113 001D 2000          LAC   ASIGN      CSIGN = ASIGN, IF ASIGN = BSIGN.
0114 001E 6008          SACL  CSIGN
0115 001F 1004          SUB   BSIGN
0116 0020 F680          BZ    OK          ELSE, CSIGN = -1.
          0021 0023
0117 0022 200C          LAC   NEGONE
0118          *      SACL  CSIGN
0119 0023 4002          OK    ZALH  AH1          SHIFT DIVIDEND TO PROTECT FROM OVERFLOW.
0120 0024 4903          ADDS  ALO
0121 0025 4B16          RPT   THREE
0122 0026 CE19          SFR
0123          *
0124 0027 4B17          RPT   FITEEN      QM = AH1:ALO / BHI, R1 = REMAINDER.
0125 0028 4706          SUBC  BHI
0126 0029 6811          SACH  R1          HIGH ACCUMULATOR RETAINS REMAINDER.
0127 002A 600F          SACL  QM
0128 002B 2F11          LAC   R1,15      (R1 * 2**15) / BHI GIVES QL, AND R2.
0129 002C 4B17          RPT   FITEEN      COMPUTES (R1 * 2**15) / BHI.
0130 002D 4706          SUBC  BHI
0131 002E 6812          SACH  R2          HIGH ACCUMULATOR RETAINS REMAINDER.
0132 002F 6010          SACL  QL
0133          *
0134 0030 3C0F          LT    QM          CORRECTION TERM = (QM * BLO) / BHI.
0135 0031 3807          MPY  BLO          COMPUTES (QM * BLO).
0136 0032 CE14          PAC
0137 0033 4B17          RPT   FITEEN      COMPUTES (QM * BLO) / BHI.
0138 0034 4706          SUBC  BHI
0139 0035 6013          SACL  CL
0140 0036 400F          ZALH  QM          QM:QL - 0:CL = CHI:CLO
0141 0037 4910          ADDS  QL
0142 0038 1013          SUB  CL
0143 0039 600B          SACL  CLO
0144 003A 680A          SACH  CHI
0145 003B 200A          LAC   CHI          DID AN OVERFLOW OCCUR?
0146 003C 4E14          AND  M1000
0147 003D F680          BZ    NOOVF       IF NOT, GOTO NOOVF.
          003E 0041
0148 003F 2015          LAC   ONE          ELSE, INCREMENT CEXP.
0149 0040 6009          SACL  CEXP
0150 0041 2001          NOOVF LAC  AEXP      COMPUTE RESULTING EXPONENT.
0151 0042 1005          SUB  BEXP
0152 0043 0009          ADD  CEXP
0153 0044 6009          SACL  CEXP
0154          *
0155          *
0156          *      NORMALIZE
0157          *
0158 0045 200A          NORMAL LAC  CHI          DOES CHI HAVE THE MSB?
0159 0046 F680          BZ    LO1
          0047 004E
0160 0048 400A          ZALH  CHI          IF YES, NORMALIZE RESULT.
0161 0049 490B          ADDS  CLO
0162 004A 4B19          RPT   TTEEN      WILL PERFORM 14 "NORMS".

```

```
0163 004B CEA2      NORM
0164 004C FF80      B      OUTPUT      GO OUTPUT RESULTS.
      004D 0057
0165 004E 400B L01  ZALH      CLO      HERE, IF CLO HAS MSB.
0166 004F F380      BLZ      NOFLOW     DID BIT SEARCH CAUSE OVERFLOW?
      0050 0055
0167 0051 4B19      RPT      TTEEN     IF NOT, NORMALIZE RESULT.
0168 0052 CEA2      NORM
0169 0053 FF80      B      OUTPUT      GO OUTPUT RESULT.
      0054 0057
0170                *
0171                *
0172                *      FINISHED WITH NORMALIZATION
0173                *
0174                *      OVERFLOW OCCURRED DURING BIT SEARCH
0175                *
0176 0055 CE06 NOFLOW RSXM          RSXM FOR LOGICAL RIGHT SHIFT.
0177 0056 CE19          SFR          PERFORM RIGHT SHIFT.
0178                *
0179                *
0180                *      TAKE CARE OF EXPONENT & NORMALIZED MANTISSA,
0181                *      THEN OUTPUT RESULTS.
0182                *
0183                *
0184 0057 680A OUTPUT SACH      CHI      SAVE NORMALIZED MANTISSA.
0185 0058 600B          SACL      CLO
0186 0059 5589          LARP      1      RESET POINTER.
0187 005A 4B16          RPT      THREE     OUTPUT RESULTS, CSIGN, CEXP, CHI, AND CLO.
0188 005B E0A0          OUT      **+,PAO
0189 005C CE1F          IDLE
NO ERRORS, NO WARNINGS
```